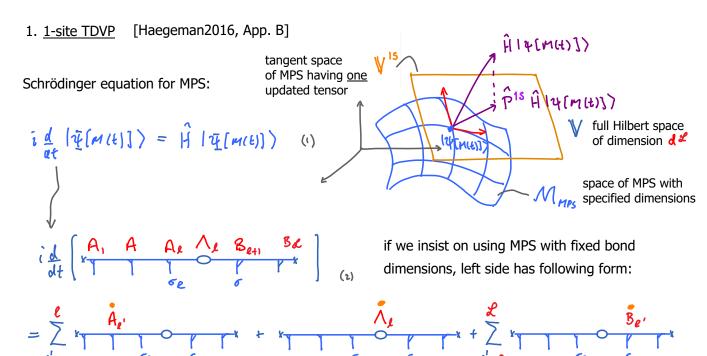
We consider time evolution using 'time-dependent variational principle' (TDVP)



Each term differs from  $|\psi(t)\rangle$  by precisely one site tensor or one bond tensor, so left side is a state in the tangent space,  $\psi$  of  $\psi$ . But right side of (1) is <u>not</u>, since since  $H \cap V(\ell)$  can have larger bond dimensions than  $V \cap V(\ell)$ .

So, project right side of (1) to 
$$V^{15}$$
:  $i \frac{d}{dt} \left| \overline{\psi} \left[ m(t) \right] \right\rangle \approx \hat{P}^{15} \left| \overline{\psi} \left[ m(t) \right] \right\rangle$  tangent space approximation

Left and right sides of (4) are structurally consistent. To see this, consider bond

Left side of (4) contains:

$$\frac{d}{dt} \frac{A_e \wedge_e \beta_{eii}}{\forall \rho} = \frac{A_e \wedge_e \beta_{eii}}{\forall \rho} + \frac{A_e \wedge_e \beta_{eii}}{\forall \rho} + \frac{A_e \wedge_e \beta_{eii}}{\forall \rho} + \frac{A_e \wedge_e \beta_{eii}}{\forall \rho}$$
(5)

Decompose: 
$$\hat{A}_{\ell} \wedge_{\ell} = A_{\ell} \wedge_{\ell} + \bar{A}_{\ell} \bar{A}_{\ell}$$
,  $A_{\ell} \bar{3}_{\ell+1} = A_{\ell}^{(1)} \bar{8}_{\ell+1} + \bar{A}^{(1)} \bar{8}_{\ell+1}$  (6)

Then we find:

Then we find:

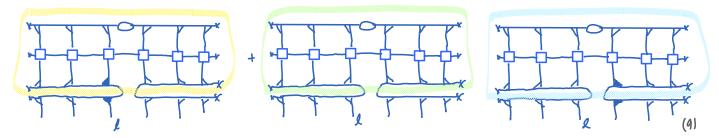
$$\frac{d}{dt} \frac{A_e \Lambda_e B_{eii}}{qt} = \frac{A_e \Lambda_e' B_{eii}}{qt} + \frac{A_e \Lambda_e' B_{eii}}{qt} + \frac{A_e \Lambda_e'' B_{eii}}{qt} + \frac{A_e \Lambda_e'' B_{eii}}{qt}$$
(7)

Right side of (4) requires tangent space projector. Consider its form (TS-I.5.25):

$$P^{1S} = \sum_{n=1}^{\infty} \frac{1}{n!} \frac{1}{n!} + \sum_{n=1}^{\infty} \frac{1}{n!} \frac{1}{n!} + \sum_{n=1}^{\infty} \frac{1}{n!} \frac{1}{n$$

$$P^{1S} = \sum_{\bar{\ell}=1}^{\ell'} \frac{1}{\ell} \sum_{\bar{\ell}+1}^{\ell'} \frac{1}{\ell'} + \sum_{\bar{\ell}=1}^{\ell'} \frac{1}{\ell'} \sum_{\bar{\ell}+1}^{\ell'} \frac{1}{\ell'} \sum_{\bar{\ell}=1}^{\ell'} \frac{1}{\ell'} \sum_{\bar{\ell}=1$$

The three terms with  $\bar{\ell} = \ell$ ,  $\ell' = \ell$ ,  $\bar{\ell} = \ell + \ell$ , applied to  $\bar{H} \setminus \bar{\Psi}(\ell)$ , yield



matching structure of (7). Thus,  $P^{ls}$ , applied to  $H(\Psi(l))$ , yields terms of precisely the right structure!

To integrate projected Schrödinger eq. (4), we write tangent space projector in the form (TS.5.26):

$$P^{IS} = \sum_{\ell=1}^{R} \frac{1}{\ell} \left( \frac{1}{\ell} \right) \left( \frac{1}{\ell} \right)$$

and write (4) as  $\frac{1}{2} + \frac{1}{2} + \frac{1}{2$ 

Right side is sum of terms, each specifying an update of one  $\psi_{\ell}^{\iota}$  or  $\psi_{\ell}^{\iota}$  on the left. Eq. (4) can be integrated one site at a time, by defining the updates through the following local Schrödinger equations:

$$i \stackrel{Ce}{\leftarrow} := \qquad \begin{array}{c} Ce \\ \downarrow \stackrel{Ce}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} \qquad \begin{array}{c} i \stackrel{Ae}{\leftarrow} \\ \vdots \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} \qquad \begin{array}{c} (12) \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array}$$

In site-canonical form, site  $\ell$  involves two terms linear in  $C_{\ell}$ :  $i \stackrel{\circ}{C}_{\ell} (t) = H_{\ell}^{15} C_{\ell} (t)$ 

Their contribution can be integrated exactly: replace  $C_{\ell}(t)$  by  $C_{\ell}(t+\tau) = e^{-iH_{\ell}^{15}\tau}C_{\ell}(t)$  (4) forward time step

In bond-canonical form, site  $\ell$  involves two terms linear in  $\Lambda_{\ell}$ :  $-i\Lambda_{\ell}(t) = H_{\ell}^{\dagger}\Lambda_{\ell}(t)$  (5)

Their contribution can be integrated exactly: replace  $\bigwedge_{\ell}$  (t) by  $\bigwedge_{\ell}$  (t- $\tau$ ) =  $e^{i H_{\ell}^b \tau} \bigwedge_{\ell}$  (t) (6)

11-TDVP Page 2

In practice,  $e^{-iH_{\ell}^{15}\tau}$  and  $e^{iH_{\ell}^{b}\tau}$  are computed by using Krylov methods.

Build a Krylov space by applying  $\mathcal{H}^{\text{IS}}_{\ell}$  multiple times to  $\mathcal{C}_{\ell}$ , set up the tridiagonal representation  $\mathcal{H}^{\text{IS}}_{\ell}$  in this basis, then compute the matrix exponential in this basis, and apply result to  $\mathcal{C}_{\ell}$ . Likewise for  $\mathcal{H}^{\text{IS}}_{\ell}$  and  $\mathcal{M}_{\ell}$ .

To successively update entire chains, alternate between site- and bond-canonical form, propagating forward or backward in time with  $H_{\ell}^{lS}$  or  $H_{\ell}^{b}$  , respectively:

1. Forward sweep, for l = 1, ..., l = 1, starting from l = 1, ..., l = 1, starting from l = 1, ..., l = 1, starting from l = 1, ..., l = 1

$$C_{\varrho}(t) \mathcal{B}_{\varrho}(t)$$

$$C_{\varrho}(t) \mathcal{B}_{\varrho}(t)$$

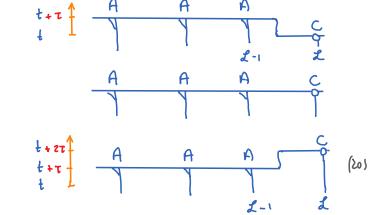
$$C_{\varrho}(t+\tau) \mathcal{B}_{\varrho}(t)$$

$$= A_{\varrho}(t+\tau) \hat{A}_{\varrho}(t+\tau) \hat{A}_{\varrho}(t+\tau) \mathcal{B}_{\varrho}(t+\tau)$$

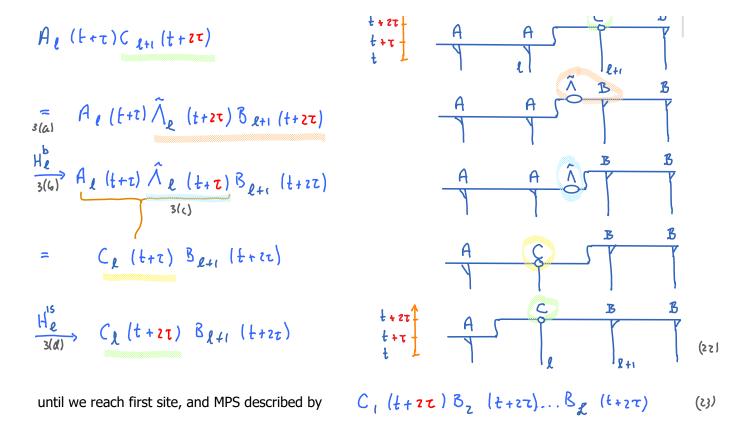
until we reach last site, and MPS described by

$$A_1(t+z)\dots A_{2n}(t+z)C_{2n}(t) \tag{19}$$

(18)



3. Backward sweep, for  $\ell = \mathcal{L} - 1, \dots, 1$ , starting from  $A_1 (t+z) \dots A_{\ell-1} (t+z) C_{\ell} (t+zz)$  (21)



The scheme described above involves 'one-site updates'. This has the (major!) drawback (as in one-site DMRG), that it is not possible to dynamically explore different symmetry sectors. To overcome this drawback, a 'two-site update' version of tangent space methods can be set up [Haegemann2016, App. C].

A systematic comparison of various MPS-based time evolution schemes has been performed in [Paeckel2019]. Conclusion: 2-site-update tangent space scheme is most accurate!

A scheme for doing 1-site TDVP while nevertheless expanding bonds, called 'controlled bond expansion (CBE), was proposed in [Li2022] (see next lecture!).

TDVP.2

The construction of tangent space  $V^{13}$  and its projector  $P^{13}$  can be generalized to n sites [Gleis2022a].

We focus on V = 2 (but general case is analogous). Define space of 2-site variations:

 $\bigvee^{25}$  = span of all states  $\bigvee^{4}$  differing from  $\bigvee^{4}$  on at most 2 neighboring sites

$$= \text{span } \left\{ \left| \vec{\Psi}' \right\rangle = r + \frac{2 \text{ sites}}{2 + 2 \text{ sites}} \right\}$$
 (1)

Recall:

Recall:
$$\frac{|\text{local 2s projector:}}{|\text{local } 2\text{s projector:}} = \frac{2\text{s}}{|\text{TS-I.4.9}|}$$
(3)

Global 2s projector  $\stackrel{\circ}{P}^{2s}$ , such that  $\stackrel{\vee}{V}^{2s} = i_{lm} \left( \stackrel{\circ}{P}^{2s} \right)$ , can be found with a Gram-Schmidt scheme analogous to our construction of  $\stackrel{\circ}{P}^{1s}$ , see [Gleis2022a]:

compare (TS-I.5.22)
$$P^{2S} := \sum_{\ell=1}^{2} P_{\ell}^{2S} + P_{\ell}^{2S} + \sum_{\ell=1}^{2} P_{\ell}^{2S} + \sum$$

$$p^{2S} = \sum_{\ell=1}^{\ell'-1} \frac{1}{\ell} \left\{ \begin{array}{c} \ell' = \ell' \\ \ell + 1 \end{array} \right\} \left\{ \begin{array}{c} \ell' = \ell' \\ \ell + 1 \end{array} \right\} \left\{ \begin{array}{c} \ell' = \ell' \\ \ell + 1 \end{array} \right\} \left\{ \begin{array}{c} \ell' = \ell' \\ \ell + 1 \end{array} \right\} \left\{ \begin{array}{c} \ell' = \ell' \\ \ell + 1 \end{array} \right\} \left\{ \begin{array}{c} \ell' = \ell' \\ \ell + 1 \end{array} \right\} \left\{ \begin{array}{c} \ell' = \ell' \\ \ell + 1 \end{array} \right\} \left\{ \begin{array}{c} \ell' = \ell' \\ \ell + 1 \end{array} \right\} \left\{ \begin{array}{c} \ell' = \ell' \\ \ell + 1 \end{array} \right\} \left\{ \begin{array}{c} \ell' = \ell' \\ \ell + 1 \end{array} \right\} \left\{ \begin{array}{c} \ell' = \ell' \\ \ell + 1 \end{array} \right\} \left\{ \begin{array}{c} \ell' = \ell' \\ \ell + 1 \end{array} \right\} \left\{ \begin{array}{c} \ell' = \ell' \\ \ell + 1 \end{array} \right\} \left\{ 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\ell + 1 \end{array} \right\} \left\{ \begin{array}{c} \ell'$$

All summands are mutually orthogonal, ensuring that  $(P^{25})^2 = P^{25}$ , and that  $P^{25} = P^{25}$ 

Alternative expression: compare (TS.5.26)

$$P^{25} = \sum_{\ell=1}^{\ell-1} P_{\ell}^{25} - \sum_{\ell=1}^{\ell-2} P_{\ell+1}^{15} = \sum_{\ell=1}^{\ell-1} \frac{1}{1} \sum_{\ell=1}^{\ell-1} \frac{1}{1}$$

This projector is used for 2-site TDVP (see TS-II.3)

## Orthogonal n-site projectors

For any given MPS |\(\vec{\psi}(M)\), full Hilbert space of chain can be decomposed into mutually orthogonal subspaces:

$$V = V_1 \otimes \cdots \otimes V_e = \bigoplus_{N=0}^{e} V^{N\perp}$$
(8)

with 
$$V^{os} := V^{os} := Span \{ | \Psi \rangle \}$$
 (9)

'irreducible' 
$$\bigvee^{NL}$$
 is complement of  $\bigvee^{(N-1)5}$  in  $\bigvee^{NS} = \bigvee^{(N-1)5} \bigoplus^{NL}$  (6)

= span of states differing from  $|\Psi\rangle$  on  $|\Psi\rangle$  contiguous sites, not expressible through subsets of  $|\Psi'\rangle$ 

Correspondingly, identity can be decomposed as:

$$1_{V} = 1_{V}^{A} = \sum_{N=0}^{R} P^{NL}$$
completeness

orthogonality

orthogonality

 $P^{\perp N}$  is defined as the projector having  $V^{N\perp}$  as image:  $I_{N}(P^{N\perp}) \approx V^{N\perp}$ (12)where

$$N \ge 1$$
:  $P^{NL} := P^{NS} (1_V - P^{(N-1)S}) = P^{NS} - P^{(N-1)S}$  (14)

since 
$$V^{(n-1)s}$$
  $\subset V^{ns} \Rightarrow im(P^{(n-1)s}) \subset im(P^{ns})$   
 $\Rightarrow P^{ns} P^{(n-1)s} = P^{(n-1)s}$ 

Consider n=1:

$$= \underbrace{\begin{array}{c} \ell' \\ \ell = 1 \end{array}}_{\ell = 1} \underbrace{\begin{array}{c} \ell' \\ \ell = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell' \\ \ell' = \ell' \end{array}}_{\ell + \ell'} + \underbrace{\begin{array}{c} \ell$$

- \*\* choose l' = L(2)

$$= \sum_{\ell=1}^{\ell} P_{\ell,\ell+1}^{bk}$$
 projects onto all 1-site variations orthogonal to  $|\Psi\rangle$ 

Consider n=2:

$$= \sum_{\ell=1}^{\ell-1} \left( \begin{array}{c} x \\ y \\ \end{array} \right) \left[ \begin{array}{c} x \\ y \\ \end{array} \right] \left[ \begin{array}{c} x \\$$

$$= \begin{array}{c} \mathcal{L}_{-1} \\ = \\ \mathcal{L}_{-1} \\ \mathcal{L}_$$

very important result! (26)

(1)

## [Haegeman2016, Sec. V & App. C]

2-site tangent space methods are analogous to 1-site methods, but use a 2-site projector. There is a conceptual difference, though: the main reason for using 2-site schemes is that they allow sectors with new quantum numbers to be introduced if the action of H requires this. However, states with different ranges of quantum numbers live in different manifolds, hence this procedure 'cannot easily be captured in a smooth evolution described using a differential equation. However, like most numerical integration schemes, the aforementioned algorithm is intrinsically discrete by choosing a time step, and it poses no problem to formulate an analogous two-site algorithm'. [Haegeman2016, Sec. V]. In other words: the tangent space approach is conceptually not as clean for the 2-site as for the 1-site scheme.

Schrödinger equation, projected onto 2-site tangent space, now takes the form

$$i \frac{d}{at} | \psi[M(t)] \rangle = \hat{\rho}^{2s} \hat{H} | \psi[M(t)] \rangle$$

$$\hat{P}^{ZS} = \sum_{\ell=1}^{2^{-1}} \frac{1}{\ell+1} - \sum_{\ell=2}^{2^{-1}} \frac{1}{\ell+1}$$

$$(2)$$

This yields [compare (1.11)]:

or 
$$\begin{cases} 2^{-1} & 4^{2} \\ 2^{-1} & 4^{2} \\ 2^{-1} & 4^{2} \end{cases}$$
  $= \begin{cases} 2^{-1} & 4^{-2} \\ 2^{-1} & 4^{-2} \\ 2^{-1} & 4^{-2} \end{cases}$   $= \begin{cases} 2^{-1} & 4^{-2} \\ 2^{-1} & 4^{-2} \\ 2^{-1} & 4^{-2} \end{cases}$   $= \begin{cases} 2^{-1} & 4^{-2} \\ 2^{-1} & 4^{-2} \\ 2^{-1} & 4^{-2} \end{cases}$ 

Right side is sum of terms, each specifying an update of one  $\psi_{\ell}^{ts}$  or  $\psi_{\ell}^{ts}$  on the left. Eq. (4) can be integrated one site at a time, by defining the updates through the following local Schrödinger equations:

$$\frac{i \dot{\psi}_{\ell}^{2S}}{1 + 1} := \frac{\psi_{\ell+1}^{2S}}{1 + 1} + \frac{i \dot{\psi}_{\ell+1}^{1S}}{\ell} := - \frac{\psi_{\ell+1}^{1S}}{1 + 1} + \frac{\psi_{\ell+1}^{1S}}{\ell} + \frac{\psi_{\ell+1}^{1S}}{1 + 1} + \frac{\psi_{\ell+1}^{1S}}{\ell} +$$

Right side is sum of terms, each linear in a factor appearing on the left. Can be integrated one site at a time:

In 2-site-canonical form, site 
$$\ell$$
 involves two terms linear in  $\Upsilon_{\ell}^{is}$ :  $i \psi_{\ell}^{is}(t) = H_{\ell}^{is} \Upsilon_{\ell}^{is}(t)$  (5)

Their contribution can be integrated exactly: replace 
$$\psi_{\ell}^{2s}(t)$$
 by  $\psi_{\ell}^{2s}(t+\tau) = e^{-i H_{\ell}^{2s} \tau} \psi_{\ell}^{2s}(t)$  (4) forward time step

In 1-site-canonical form, site 
$$\ell$$
+1 involves two terms linear in  $\Psi_{\ell+1}^{15}$ :  $-i \psi_{\ell+1}^{15}(t) = H_{\ell+1}^{15}(t)$ 

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Their contribution can be integrated exactly: replace 
$$\psi_{\ell+1}^{(5)}(t)$$
 by  $\psi_{\ell+1}^{(5)}(t-\tau) = e^{iH_{\ell+1}^{(5)}\tau} \psi_{\ell+1}^{(5)}(t)$  backward(!) time step

To successively update entire chains, alternate between 2-site- and 1-site-canonical form, propagating forward or backward in time with  $H_{\ell}^{\text{rs}}$  or  $H_{\ell}^{\text{rs}}$ , respectively (analogously to 1-site scheme).

A systematic comparison of various MPS-based time evolution schemes has been performed in [Paeckel2019]. Conclusion: 2-site-update tangent space scheme is most accurate!