

II: Numerical linear algebra

Within linear algebra, we perform arithmetic operations on finite precision arithmetic.

In the following we consider a finite dim. vector space over real numbers $V_{\mathbb{R}}^n$ ($n \in \mathbb{N}$).

The complex case is obtained from $n \rightarrow 2n$.
Available, elementary arithmetic operations:

- $+, -$: n -digit arguments typically $\sim O(n)$ complexity
- $\cdot, /$: n -digit arguments typically $\sim O(n^2)$ complexity

Take home message:

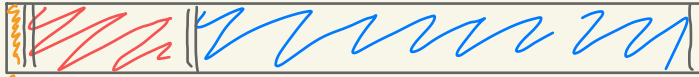
Complexity is estimated w.r.t. # multiplications/divisions!

II.1: Finite precision arithmetics

In finite precision arithmetics there always is a discretization Δ in the representation.

Example: Double precision (64 bit float)

63 62 - - - 52 - - - - - 0 bits



1 sign bit $n_e = 11$ exponent bits $n_f = 52$ fraction bits

Representation of real number x :

$$x = (-1)^{\text{sign}} \cdot \left(1 + \sum_{n=1}^{n_f} x_{n_f-n} \cdot 2^{-n} \right) \cdot 2^{e-1023}$$

with $e = \sum_{n=0}^{n_e-1} x_{n_f+n_e-n} 2^n$ the exponent.

Thus here: $\Delta = 2^{-52} \cdot 2^{e-1023}$

Define relative precision: $\delta = 2^{-n_f}$

Thus here: $\delta = 2^{-52} = (1024)^{-5} \cdot \frac{1}{4} \approx 10^{-16}$

Consequences:

(i) $x \in \mathbb{R}$ only represented modulo δ :

$$\tilde{x} = x(1 + \delta_x), \quad |\delta_x| \leq \delta/2$$

(ii) $z = x + y$ with rounding errors δ_x, δ_y :

$$\tilde{z} = \tilde{x} + \tilde{y} = z + \underbrace{(x\delta_x + y\delta_y)}_{z\delta_z}, \quad |\delta_z| \leq \delta/2$$

(iii) $z = x - y$ with rounding errors δ_x, δ_y :

$$\tilde{z} = \tilde{x} - \tilde{y} = z + (x\delta_x - y\delta_y) = z \left(1 + \underbrace{\frac{x\delta_x - y\delta_y}{z}}_{\delta_z} \right)$$

Note: δ_z here not bounded but depends on z ! $z \rightarrow 0$, then δ_z can explode!

called: Catastrophic cancellation

(iv) $z = x \cdot y$ with rounding errors δ_x, δ_y :

$$\tilde{z} = \tilde{x} \cdot \tilde{y} = x \cdot y (1 + \delta_x + \delta_y + \delta_x \cdot \delta_y) = z (1 + \delta_z)$$

note that if $\tilde{x} \cdot \tilde{y} < \delta$, then $\delta_z > \delta$ possible.

To see this write

$$\tilde{x} = (1 + r_x) 2^{e_x - 1025}, \quad \tilde{y} = (1 + r_y) 2^{e_y - 1025}, \quad |r_{x,y}| < 1$$

$$\Rightarrow \tilde{x} \cdot \tilde{y} = (1 + r_x)(1 + r_y) 2^{e_x + e_y - 2046}$$

consider $r_x r_y 2^{e_x + e_y - 2046}$

$$= \sum_{n,m=0}^{n_f} \left(X_{n_f-n} Y_{n_f-m} 2^{-(n+m)} \right)$$

$$\text{and thus: } |\delta_z| = \sum_{n+m > n_f} X_{n_f-n} Y_{n_f-m} 2^{-(n+m)}$$

\Rightarrow multiplication suppresses relevant bits

(3)

(v) $z = x/y$ similar to $x \cdot y$ but here errors are magnified (can be worked out similar to (iv))

Keep these rounding errors in mind as they can drastically impact outcome!

Example:

$$(i) \quad f(x) = \frac{x - \sin(x)}{1 - \cos(x)} \equiv y$$

Taylor expand $\sin(x)$ & $\cos(x)$ & evaluate at finite precision arithmetics for $x \ll 1$

$$\tilde{y} = \frac{x(\delta_x - \delta_x)}{x^2(1 - \delta_x)^2}$$

numerator evaluates with precision δ but denominator with max precision $\sqrt{\delta}$!

$\Rightarrow x < \sqrt{\delta}$ then y undefined

$$(ii) \quad \text{Var}(X) = \langle (X - \langle X \rangle)^2 \rangle$$

evaluating $\text{Var}(X) = \langle X^2 \rangle - \langle X \rangle^2$

limits precision to $\delta \text{Var}(X) \sim \sqrt{\delta}$ (catastrophic cancellation)

(iii) $z = x^2 - y^2$

evaluating $z = (x^2) - (y^2)$ limits precision to

$\sqrt{\delta}$. But evaluating $z = (x+y)(x-y)$ evaluates

to precision δ !

II.2 Matrix-Matrix multiplication

In the following, we consider for simplicity quadr. matrices $\in \mathbb{V}_{\mathbb{R}}^{m \times m}$ with $m = 2^n$, $n \in \mathbb{N}$.

Divide and conquer

Decompose $\underline{\underline{C}} = \underline{\underline{A}} \cdot \underline{\underline{B}}$ with $\underline{\underline{A}}, \underline{\underline{B}} \in \mathbb{V}_{\mathbb{R}}^{m \times m}$ into

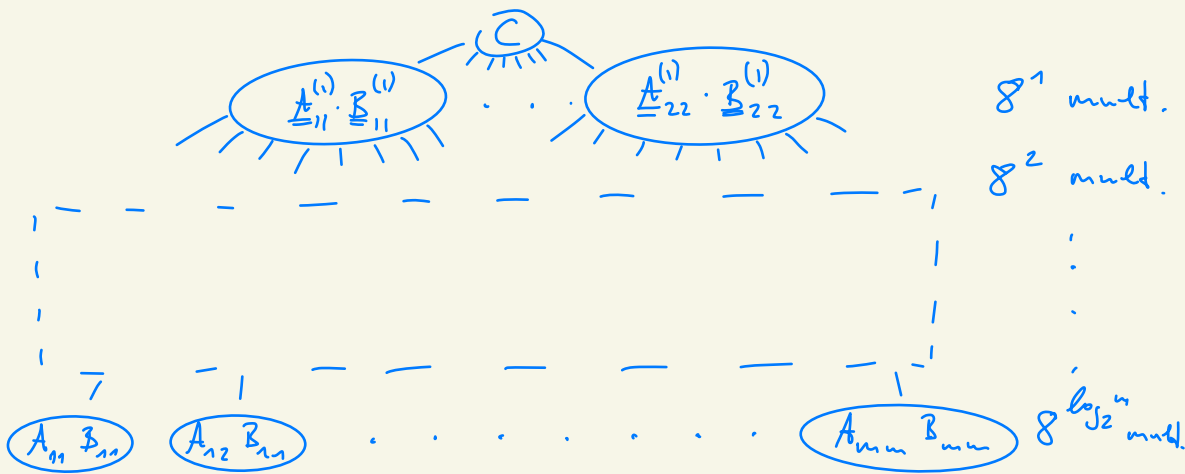
$$\begin{pmatrix} \underline{\underline{C}}_{11} & \underline{\underline{C}}_{12} \\ \underline{\underline{C}}_{21} & \underline{\underline{C}}_{22} \end{pmatrix} = \begin{pmatrix} \underline{\underline{A}}_{11} & \underline{\underline{A}}_{12} \\ \underline{\underline{A}}_{21} & \underline{\underline{A}}_{22} \end{pmatrix} \cdot \begin{pmatrix} \underline{\underline{B}}_{11} & \underline{\underline{B}}_{12} \\ \underline{\underline{B}}_{21} & \underline{\underline{B}}_{22} \end{pmatrix}$$

with $\underline{\underline{A}}_{ij}, \underline{\underline{B}}_{ij}, \underline{\underline{C}}_{ij} \in \mathbb{V}_{\mathbb{R}}^{m/2 \times m/2}$ and

$$\begin{aligned}
 \underline{\underline{C}}_{11} &= \underline{\underline{A}}_{11} \cdot \underline{\underline{B}}_{11} + \underline{\underline{A}}_{12} \cdot \underline{\underline{B}}_{21} \\
 \underline{\underline{C}}_{12} &= \underline{\underline{A}}_{11} \cdot \underline{\underline{B}}_{21} + \underline{\underline{A}}_{12} \cdot \underline{\underline{B}}_{22} \\
 \underline{\underline{C}}_{21} &= \underline{\underline{A}}_{21} \cdot \underline{\underline{B}}_{11} + \underline{\underline{A}}_{22} \cdot \underline{\underline{B}}_{21} \\
 \underline{\underline{C}}_{22} &= \underline{\underline{A}}_{21} \cdot \underline{\underline{B}}_{12} + \underline{\underline{A}}_{22} \cdot \underline{\underline{B}}_{22}
 \end{aligned}
 \left. \vphantom{\begin{aligned} \underline{\underline{C}}_{11} \\ \underline{\underline{C}}_{12} \\ \underline{\underline{C}}_{21} \\ \underline{\underline{C}}_{22} \end{aligned}} \right\} 8 \text{ M-M-operations}$$

For each M-M-operation $\underline{\underline{A}}_{ij} \cdot \underline{\underline{B}}_{kl}$, $i, j, k, l \in \{1, 2\}$ repeat scheme. After $\log_2 m$ recursions we arrive at 8 simple multiplications.

The total number of multiplications can be deduced from drawing iteration tree:



$$\Rightarrow \# \text{ multiplications} = 8^{\log_2 m} = (2^3)^{\log_2 m} = m^3$$

Even though that is not surprising, it tells

us two important things:

(i) $M-M$ products can be parallelized systematically

(ii) The counting technique suggests investigating the decomposition for more efficient scheme

Indeed, the 8 equations are highly symmetric!

Strassen's algorithm (1969)

There are various ways to prove Strassen's algorithm. Let us begin by introducing a basis for $V_{\mathbb{R}}^{2 \times 2}$.

$$\begin{aligned} \underline{e}^1 = |e^1\rangle\langle e^1| &= \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}, & \underline{e}^2 = |e^2\rangle\langle e^2| &= \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix} & \text{with } |e^1\rangle &= \begin{pmatrix} 1 \\ 0 \end{pmatrix} \\ \underline{e}^3 = |e^2\rangle\langle e^1| &= \begin{pmatrix} 0 & 0 \\ 1 & 0 \end{pmatrix}, & \underline{e}^4 = |e^2\rangle\langle e^2| &= \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix} & |e^2\rangle &= \begin{pmatrix} 0 \\ 1 \end{pmatrix} \end{aligned}$$

They obey the algebra:

$$\left. \begin{array}{cccc} \underline{e}^1 \cdot \underline{e}^1 = \underline{e}^1 & \underline{e}^2 \cdot \underline{e}^1 = 0 & \underline{e}^3 \cdot \underline{e}^1 = \underline{e}^3 & \underline{e}^4 \cdot \underline{e}^1 = 0 \\ \underline{e}^1 \cdot \underline{e}^2 = \underline{e}^2 & \underline{e}^2 \cdot \underline{e}^2 = 0 & \underline{e}^3 \cdot \underline{e}^2 = \underline{e}^4 & \underline{e}^4 \cdot \underline{e}^2 = 0 \\ \underline{e}^1 \cdot \underline{e}^3 = 0 & \underline{e}^2 \cdot \underline{e}^3 = \underline{e}^1 & \underline{e}^3 \cdot \underline{e}^3 = 0 & \underline{e}^4 \cdot \underline{e}^3 = \underline{e}^3 \\ \underline{e}^1 \cdot \underline{e}^4 = 0 & \underline{e}^2 \cdot \underline{e}^4 = \underline{e}^2 & \underline{e}^3 \cdot \underline{e}^4 = 0 & \underline{e}^4 \cdot \underline{e}^4 = \underline{e}^4 \end{array} \right\} (\ast)$$

Now, we can write for any $\underline{A} \in \mathbb{V}_{\mathbb{R}}^{m \times m}$

$$\begin{aligned}\underline{A} &= \underline{A}_{11} \otimes \underline{e}^1 + \underline{A}_{12} \otimes \underline{e}^2 + \underline{A}_{21} \otimes \underline{e}^3 + \underline{A}_{22} \otimes \underline{e}^4 \\ &= \underline{A}_1 \otimes \underline{e}^1 + \underline{A}_2 \otimes \underline{e}^2 + \underline{A}_3 \otimes \underline{e}^3 + \underline{A}_4 \otimes \underline{e}^4\end{aligned}$$

Note that

$$\begin{aligned}\underline{A} + \lambda \underline{B} &= (\underline{A}_1 + \lambda \underline{B}_1) \otimes \underline{e}^1 + (\underline{A}_2 + \lambda \underline{B}_2) \otimes \underline{e}^2 + \dots \\ &= \sum_{j=1}^4 (\underline{A}_j + \lambda \underline{B}_j) \otimes \underline{e}^j\end{aligned}$$

We can thus interpret any $\underline{A} \in \mathbb{V}_{\mathbb{R}}^{m \times m}$ as linear combination of states formed from $\text{map}^n \otimes^n$

$$\text{via: } \otimes : \mathbb{V}_{\mathbb{R}}^{m/2 \times m/2} \times \mathbb{V}_{\mathbb{R}}^{2 \times 2} \rightarrow \mathbb{V}_{\mathbb{R}}^{m \times m}$$

We now expand $\underline{C} = \underline{A} \cdot \underline{B}$ using different basis sets $\{\underline{b}^i\}, \{\tilde{\underline{b}}^j\}$

$$\begin{aligned}\underline{C} &= \sum_{i,j=1}^4 (\underline{A}_i \cdot \underline{B}_j) \otimes (\underline{b}^i \cdot \tilde{\underline{b}}^j) \\ &= (\underline{A}_1 \cdot \underline{B}_1) \otimes (\underline{b}^1 \cdot \tilde{\underline{b}}^1) + (\underline{A}_1 \cdot \underline{B}_2) \otimes (\underline{b}^1 \cdot \tilde{\underline{b}}^2) + (\underline{A}_1 \cdot \underline{B}_3) \otimes (\underline{b}^1 \cdot \tilde{\underline{b}}^3) + (\underline{A}_1 \cdot \underline{B}_4) \otimes (\underline{b}^1 \cdot \tilde{\underline{b}}^4) \\ &\quad + (\underline{A}_2 \cdot \underline{B}_1) \otimes (\underline{b}^2 \cdot \tilde{\underline{b}}^1) + \dots \\ &\quad \vdots \\ &\quad + (\underline{A}_4 \cdot \underline{B}_1) \otimes (\underline{b}^4 \cdot \tilde{\underline{b}}^1) + \dots + (\underline{A}_4 \cdot \underline{B}_4) \otimes (\underline{b}^4 \cdot \tilde{\underline{b}}^4)\end{aligned}$$

If we would choose for the canonical basis sets $\underline{b}^i = \underline{\tilde{b}}^j = \underline{e}^j$, then using (*) we would arrive at

$$\begin{aligned} \underline{c} = & (\underline{A}_1 \cdot \underline{B}_1) \otimes \underline{e}^1 + (\underline{A}_1 \cdot \underline{B}_2) \otimes \underline{e}^2 \\ & + (\underline{A}_2 \cdot \underline{B}_3) \otimes \underline{e}^1 + (\underline{A}_2 \cdot \underline{B}_4) \otimes \underline{e}^2 \\ & + (\underline{A}_3 \cdot \underline{B}_1) \otimes \underline{e}^3 + (\underline{A}_3 \cdot \underline{B}_2) \otimes \underline{e}^4 \\ & + (\underline{A}_4 \cdot \underline{B}_3) \otimes \underline{e}^3 + (\underline{A}_4 \cdot \underline{B}_4) \otimes \underline{e}^4 \end{aligned}$$

which is just the result from the divide-and-conquer decomposition.

Thus, the question reduces to:

Can we find basis sets $\underline{b}^i, \underline{\tilde{b}}^j$ such that we need less multiplications?

Let us introduce: $\underline{D} = \begin{pmatrix} 0 & -1 \\ 1 & -1 \end{pmatrix} = \underline{e}^3 - (\underline{e}^2 + \underline{e}^4)$.

Note that \underline{D} is rotation matrix & $\underline{D}^3 = \underline{1} \Rightarrow \underline{D}^2 = \underline{D}^{-1}$

Then choose $\underline{X} = \underline{e}^2$ (important: $\underline{D}\underline{X} \neq \underline{X}$)

we have:

$$\begin{aligned} \underline{D}^{-1} = \underline{D}^2 &= \underline{e}^2 - \underline{e}^1 - \underline{e}^3 \\ \underline{D} \underline{X} \underline{D} &= \underline{e}^3 - \underline{e}^4 \end{aligned}$$

and:

$$\underline{\underline{X}} \underline{\underline{1}} \underline{\underline{X}} = \underline{\underline{e}}^2 \cdot \underline{\underline{e}}^2 = 0$$

$$\underline{\underline{X}} \underline{\underline{D}} \underline{\underline{X}} = \underline{\underline{e}}^2 (\underline{\underline{e}}^3 - \underline{\underline{e}}^2 - \underline{\underline{e}}^4) \underline{\underline{e}}^2 = (\underline{\underline{e}}^1 - \underline{\underline{e}}^2) \underline{\underline{e}}^2 = \underline{\underline{e}}^2 = \underline{\underline{X}}$$

$$\underline{\underline{X}} \underline{\underline{D}}^2 \underline{\underline{X}} = (\underline{\underline{e}}^1 - \underline{\underline{e}}^2) \underline{\underline{e}}^4 = -\underline{\underline{e}}^2 = -\underline{\underline{X}}$$

it can be shown that the two sets are a basis of $V_{\mathbb{R}}^{2 \times 2}$:

$$\mathcal{B}_1 = \{ \underline{\underline{D}}, \underline{\underline{X}}, \underline{\underline{D}}^{-1} \underline{\underline{X}} \underline{\underline{D}}, \underline{\underline{D}} \underline{\underline{X}} \underline{\underline{D}}^{-1} \} = \{ \underline{\underline{b}}^1, \underline{\underline{b}}^2, \underline{\underline{b}}^3, \underline{\underline{b}}^4 \}$$

$$\mathcal{B}_2 = \{ \underline{\underline{D}}^{-1}, \underline{\underline{X}}, \underline{\underline{D}}^{-1} \underline{\underline{X}} \underline{\underline{D}}, \underline{\underline{D}} \underline{\underline{X}} \underline{\underline{D}}^{-1} \} = \{ \underline{\underline{b}}^1, \underline{\underline{b}}^2, \underline{\underline{b}}^3, \underline{\underline{b}}^4 \}$$

Now we compute all products of $\underline{\underline{b}}^i \in \mathcal{B}_1$ & $\underline{\underline{b}}^j \in \mathcal{B}_2$

$\underline{\underline{b}}^i \backslash \underline{\underline{b}}^j$	$\underline{\underline{D}}^{-1}$	$\underline{\underline{X}}$	$\underline{\underline{D}}^{-1} \underline{\underline{X}} \underline{\underline{D}}$	$\underline{\underline{D}} \underline{\underline{X}} \underline{\underline{D}}^{-1}$
$\underline{\underline{D}}$	$\underline{\underline{1}}$	$\underline{\underline{D}} \underline{\underline{X}}$	$\underline{\underline{X}} \underline{\underline{D}}$	$\underline{\underline{D}}^{-1} \underline{\underline{X}} \underline{\underline{D}}^{-1}$
$\underline{\underline{X}}$	$\underline{\underline{X}} \underline{\underline{D}}^{-1}$	0	$-\underline{\underline{X}} \underline{\underline{D}}$	$\underline{\underline{X}} \underline{\underline{D}}^{-1}$
$\underline{\underline{D}}^{-1} \underline{\underline{X}} \underline{\underline{D}}$	$\underline{\underline{D}}^{-1} \underline{\underline{X}}$	$\underline{\underline{D}}^{-1} \underline{\underline{X}}$	0	$-\underline{\underline{D}}^{-1} \underline{\underline{X}} \underline{\underline{D}}^{-1}$
$\underline{\underline{D}} \underline{\underline{X}} \underline{\underline{D}}^{-1}$	$\underline{\underline{D}} \underline{\underline{X}} \underline{\underline{D}}$	$-\underline{\underline{D}} \underline{\underline{X}}$	$\underline{\underline{D}} \underline{\underline{X}} \underline{\underline{D}}$	0

Note that only seven products occur!

Thus we can expand $\underline{\underline{C}} = \underline{\underline{A}} \cdot \underline{\underline{B}}$ in this basis using

$$\underline{\underline{b}}^1 \cdot \underline{\underline{b}}^1 = \underline{\underline{1}} = \underline{\underline{e}}^1 + \underline{\underline{e}}^4$$

$$\underline{\underline{b}}^1 \cdot \underline{\underline{b}}^2 = \underline{\underline{D}} \underline{\underline{X}} = \underline{\underline{e}}^4$$

$$\underline{\underline{b}}^2 \cdot \underline{\underline{b}}^1 = \underline{\underline{X}} \underline{\underline{D}}^{-1} = -\underline{\underline{e}}^1$$

$$\underline{\underline{b}}^1 \cdot \underline{\underline{b}}^4 = \underline{\underline{D}}^{-1} \underline{\underline{X}} \underline{\underline{D}}^{-1} = \underline{\underline{e}}^1 + \underline{\underline{e}}^3$$

$$\underline{\underline{b}}^1 \cdot \underline{\underline{b}}^3 = \underline{\underline{X}} \underline{\underline{D}} = \underline{\underline{e}}^1 - \underline{\underline{e}}^2$$

$$\underline{\underline{b}}^3 \cdot \underline{\underline{b}}^1 = \underline{\underline{D}}^{-1} \underline{\underline{X}} = -\underline{\underline{e}}^2 - \underline{\underline{e}}^4$$

$$\underline{\underline{b}}^4 \cdot \underline{\underline{b}}^1 = \underline{\underline{D}} \underline{\underline{X}} \underline{\underline{D}} = \underline{\underline{e}}^3 - \underline{\underline{e}}^4$$

such that

$$\underline{\underline{C}} = (\underline{\underline{A}}_1 \cdot \underline{\underline{B}}_1) \otimes (\underline{\underline{e}}^1 + \underline{\underline{e}}^4)$$

$$+ (\underline{\underline{A}}_1 - \underline{\underline{A}}_4) \cdot \underline{\underline{B}}_2 \otimes \underline{\underline{e}}^4$$

$$+ (\underline{\underline{A}}_1 - \underline{\underline{A}}_2) \cdot \underline{\underline{B}}_3 \otimes (\underline{\underline{e}}^1 - \underline{\underline{e}}^2)$$

$$- \underline{\underline{A}}_2 (\underline{\underline{B}}_1 + \underline{\underline{B}}_4) \otimes \underline{\underline{e}}^1$$

(*)

$$- \underline{\underline{A}}_3 (\underline{\underline{B}}_1 + \underline{\underline{B}}_2) \otimes (\underline{\underline{e}}^2 + \underline{\underline{e}}^4)$$

$$+ (\underline{\underline{A}}_1 - \underline{\underline{A}}_3) \underline{\underline{B}}_4 \otimes (\underline{\underline{e}}^1 + \underline{\underline{e}}^3)$$

$$+ \underline{\underline{A}}_4 (\underline{\underline{B}}_1 + \underline{\underline{B}}_3) \otimes (\underline{\underline{e}}^3 - \underline{\underline{e}}^4)$$

Now we only need to represent $\underline{\underline{A}}_i$ & $\underline{\underline{B}}_i$ in the old basis $\underline{\underline{e}}^i$:

$$\underline{\underline{b}}^1 = \underline{\underline{e}}^3 - \underline{\underline{e}}^2 - \underline{\underline{e}}^4$$

$$\underline{\underline{b}}^2 = \underline{\underline{e}}^2$$

$$\underline{\underline{b}}^3 = \underline{\underline{e}}^2 - \underline{\underline{e}}^1 - \underline{\underline{e}}^3$$

$$\underline{\underline{b}}^4 = \underline{\underline{e}}^2$$

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$$\underline{b}^3 = -\underline{e}^1 + \underline{e}^2 - \underline{e}^3 + \underline{e}^4$$

$$\underline{b}^{23} = -\underline{e}^1 + \underline{e}^2 - \underline{e}^3 + \underline{e}^4$$

$$\underline{b}^4 = -\underline{e}^3$$

$$\underline{b}^{14} = -\underline{e}^3$$

$$\Rightarrow \underline{A} = \underline{A}_1 \otimes (\underline{e}^3 - \underline{e}^2 - \underline{e}^4)$$

$$\Rightarrow \underline{B} = \underline{B}_1 \otimes (\underline{e}^2 - \underline{e}^1 - \underline{e}^3)$$

$$+ \underline{A}_2 \otimes \underline{e}^2$$

$$+ \underline{B}_2 \otimes \underline{e}^2$$

$$+ \underline{A}_3 \otimes (-\underline{e}^1 + \underline{e}^2 - \underline{e}^3 + \underline{e}^4)$$

$$+ \underline{B}_3 \otimes (-\underline{e}^1 + \underline{e}^2 - \underline{e}^3 + \underline{e}^4)$$

$$- \underline{A}_4 \otimes \underline{e}^3$$

$$- \underline{B}_4 \otimes \underline{e}^3$$

$$\stackrel{!}{=} \underline{A}_{11} \otimes \underline{e}^1 + \underline{A}_{12} \otimes \underline{e}^2$$

$$\stackrel{!}{=} \underline{B}_{11} \otimes \underline{e}^1 + \underline{B}_{12} \otimes \underline{e}^2$$

$$+ \underline{A}_{21} \otimes \underline{e}^3 + \underline{A}_{22} \otimes \underline{e}^4$$

$$+ \underline{B}_{21} \otimes \underline{e}^3 + \underline{B}_{22} \otimes \underline{e}^4$$

solving for \underline{A}_i ($\underline{A}_{11}, \dots, \underline{A}_{22}$) and \underline{B}_i ($\underline{B}_{11}, \dots, \underline{B}_{22}$) and inserting into (**), after some further algebra one arrives at

$$\underline{c} = (\underline{M}_I + \underline{M}_{IV} - \underline{M}_{V} + \underline{M}_{VI}) \otimes \underline{e}^1$$

$$+ (\underline{M}_{III} + \underline{M}_{VI}) \otimes \underline{e}^2 + (\underline{M}_{II} + \underline{M}_{IV}) \otimes \underline{e}^3$$

$$+ (\underline{M}_I + \underline{M}_{III} - \underline{M}_{II} + \underline{M}_{VI}) \otimes \underline{e}^4$$

with

$$\underline{M}_I = (\underline{A}_{11} + \underline{A}_{22})(\underline{B}_{11} + \underline{B}_{22}), \quad \underline{M}_{II} = (\underline{A}_{11} + \underline{A}_{22}) \underline{B}_{11}$$

$$\underline{M}_{III} = \underline{A}_{11}(\underline{B}_{12} - \underline{B}_{22}), \quad \underline{M}_{IV} = \underline{A}_{22}(-\underline{B}_{11} + \underline{B}_{21}), \quad \underline{M}_{V} = (\underline{A}_{11} + \underline{A}_{12}) \underline{B}_{22}$$

$$\underline{M}_{VI} = (-\underline{A}_{11} + \underline{A}_{21})(\underline{B}_{11} + \underline{B}_{12}), \quad \underline{M}_{VII} = (\underline{A}_{12} - \underline{A}_{22})(\underline{B}_{21} + \underline{B}_{22})$$

We can analyze the complexity by drawing the same tree of iterations but now with only 7 multiplications, i.e. 7 branches emerging from each node.

$$\Rightarrow \# \text{ multiplications} = 7^{\log_2 n} = n^{\log_2 7} \quad !!!$$