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# Sheet 1: Linear Algebra 

Released: 04/21/23; Submit until: 05/05/23 (20 Points)

This sheet is about a very small introduction into elementary concepts of C++ and how to use them to implement fast linear algebra operations using Intel's Math Kernel Library (MKL). In order to work through the problems, checkout the git repository nqp-exercises provided on the lecture's homepage and work with the code templates. While the header file la_wrapper.h, which contains an interface to some linear algebra methods of the MKL, is complete, in each code template there are files which are erroneous and need to be corrected. In a Jupyter-Hub session, use the provided Makefile as well as a proper make.inc to compile the corrected code templates.

Problem 1 Implementing a matrix class (5 Points)
Checkout the code template blas_wrapper, which contains an elementary implementation of a wrapper of the widely used linear algebra MKL.
(1.a) (3P) Correct the code template blas_wrapper such that you can compile it using make build. As a consistency check, you can use the default implementation of the main-routine and check the output when executing the compiled binary.
(1.b) (2P) You now have a rudimentary matrix class that implements fast linear algebra operations using Intel's MKL. Use this class and write a main-routine which performs a scaling analysis of the runtime needed to copy matrices with dimension $m \times m$ for $m \in \mathbb{N}$ and plot the runtime as a function of $m$. Extract the exponent $\alpha$ determining the dominating scaling of the runtime $t \sim m^{\alpha}$.

Problem 2 Implementing fast matrix contractions (10 Points)
For this exercise, checkout the code template gemm_wrapper, which provides an elementary implementation of the fast xgemm operations. Here, the $x$ denotes the fundamental data type, i.e., single, double, complex single, or complex double. The acronym gemm abbreviates GEneral Matrix-Matrix multiplications and this convention carries over for other provided operations, for instance, GEneral Matrix-Vector multiplications (checkout Intel's developer guide for a rather complete documentation about the blas/lapack interface).
(2.a) (5P) Proceed as in problem (1) and correct the code template. Note how la_operations.h now also contains a multiplication operations, which is compiled into an object file la_objects.o that now implements a general matrix-matrix multiplication. Write a test case, which tests the implemented xgemm-functionality and returns success (return code 0 ) or failure (return code 1) on exit, depending on whether a successful matrix-matrix multiplication has been performed. Why is such a test case useful?
(2.b) (5P) Write a trivial version of a matrix-matrix multiplication by implementing the calculations of the elements

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\begin{equation*}
C_{i j}=\sum_{k} A_{i k} B_{k j} \tag{1}
\end{equation*}
$$

of the result of a matrix-matrix product $\mathbf{C}=\mathbf{A} \cdot \mathbf{B}$, explicitly. Here, we consider $\mathbf{A}, \mathbf{B}, \mathbf{C} \in$ $\mathbb{V}_{\mathbb{R}}^{m \times m}$ for some $m \in \mathbb{N}$. Perform a runtime analysis comparing your naive implementation with xgemm calls, investigate the dependency of the runtime on the matrix dimension $m$ and extract the exponent $\beta$ determining the dominating scaling of the runtime $t \sim m^{\beta}$. Interpret you result.

## Problem 3 Syntactic sugar for matrix contractions (5 Points)

For this exercise, checkout the code template expr_templates, which provides an elementary implementation of expression templates to overload the multiplication operator. In problem (2) we introduced an operator overload to the multiplication operator *, allowing for expression such as $\mathrm{C}=\mathrm{A} * \mathrm{~B}$ in $\mathrm{C}^{+}$. However, that implementation also required an intermediate copy operation, which is necessary because $*$ is a binary operator and the result of the matrix-matrix multiplication has to get through a temporary return value. This unfortunate fact can be avoided by delayed evaluation, which in $\mathrm{C}^{++}$can be implemented using expression templates.
(3.a) (3P) Proceed as in problem (1) and (2) and correct the code template. In particular note how the binary matrix-matrix multiplication operator $*$ is mapped to the unary assignment operator $=$, which assigns the result of the operation to an instance of LAMatrix without additional copy-operations.
(3.b) (2P) Perform a runtime analysis and extract the speed-up as a function of the matrixdimension $m$ obtained, using the xgemm-impementations of the *-operator from problem (2) and (3).

