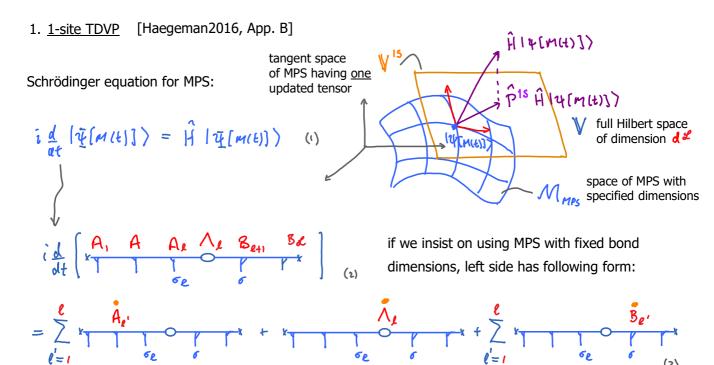
We consider time evolution using 'time-dependent variational principle' (TDVP)



Each term differs from $|\Psi(t)\rangle$ by precisely one site tensor or on bond tensor, so left side is a state in the tangent space, $\sqrt{15}$ of $\sqrt{\psi(1)}$. But right side of (1) is <u>not</u>, since since $H (\Psi(\ell))$ can have larger bond dimensions than $\Psi(\ell)$.

So, project right side of (1) to
$$V^{15}$$
: $i \frac{d}{dt} \left(\frac{1}{4} \left[m(t) \right] \right) \stackrel{\sim}{=} P^{15} \left[\frac{1}{4} \left[m(t) \right] \right)$ (4) tangent space approximation

Left and right sides of (4) are structurally consistent. To see this, consider bond

Left side of (4) contains:

$$\frac{d}{dt} \frac{A_e \wedge_e \beta_{e_{i_1}}}{\forall p} = \frac{A_e \wedge_e \beta_{e_{i_1}}}{\forall p} + \frac{A_e \wedge_e \beta_{e_{i_1}}}{\forall p} + \frac{A_e \wedge_e \beta_{e_{i_1}}}{\forall p}$$
(5)

Decompose:
$$\hat{A}_{\ell} = A_{\ell} \Lambda_{\ell}^{\prime} + \bar{A}_{\ell} \bar{\Lambda}_{\ell}^{\prime}$$
, $\mathcal{F}_{2+1} = \Lambda_{\ell}^{\prime\prime} \mathcal{F}_{\ell+1} + \bar{\Lambda}^{\prime\prime} \bar{\mathcal{F}}_{\ell+1}$ (6)

Then we find:

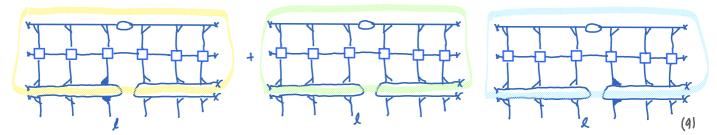
nen we find:
$$\frac{d}{dt} \frac{A_e \Lambda_e \, \delta_{e_{41}}}{dt} = \frac{\overline{A_e \, \Lambda_e \, \delta_{e_{41}}}}{T} + \frac{A_e \, \overline{\Lambda_e \, \delta_{$$

Right side of (4) requires tangent space projector. Consider its form (5.25):

$$P^{1S} = \sum_{n=1}^{\infty} \frac{1}{n} \frac{1}{n}$$

$$e^{is} = \sum_{\bar{\ell}=1}^{e'} \frac{1}{\ell} \sum_{\bar{\ell}+i}^{e'} \frac{1$$

The three terms with $\bar{\ell} = \ell$, $\ell' = \ell$, $\bar{\ell} = \ell + \ell$, applied to , $\hat{H} \setminus \bar{\Psi}(4)$, yield



matching structure of (7). Thus, P^{15} , applied to $H(\Psi(1))$, yields terms of precisely the right structure!

To integrate projected Schrödinger eq. (4), we write tangent space projector in the form (5.26):

$$P^{IS} = \sum_{\ell=1}^{R} \frac{1}{\ell} \left(\frac{1}{\ell} \right) \left(\frac{1}{\ell} \right)$$

and write (4) as $\frac{1}{2} = \frac{1}{2} = \frac{1}{2$

Right side is sum of terms, each specifying an update of one ψ_{ℓ}^{ls} or ψ_{ℓ}^{ls} on the left. Eq. (4) can be integrated one site at a time, by defining the updates through the following local Schrödinger equations:

$$i \stackrel{Ce}{\leftarrow} := \qquad \begin{array}{c} Ce \\ \downarrow \stackrel{Ce}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \downarrow \stackrel{Ae}{\leftarrow} \\ \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \longrightarrow \qquad \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \longrightarrow \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \longrightarrow \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \longrightarrow \end{array} := \qquad \begin{array}{c} \stackrel{Ae}{\leftarrow} \\ \longrightarrow \qquad \begin{array}{c$$

In site-canonical form, site ℓ involves two terms linear in C_{ℓ} : $i \stackrel{\circ}{C}_{\ell} (t) = H_{\ell}^{15} C_{\ell} (t)$

Their contribution can be integrated exactly: replace $C_{\ell}(t)$ by $C_{\ell}(t+\tau) = e^{-i\frac{H_{\ell}^{15}\tau}{2}}C_{\ell}(t)$ forward time step

In bond-canonical form, site ℓ involves two terms linear in Λ_{ℓ} : $i \Lambda_{\ell}(t) = -H_{\ell}^{b} \Lambda_{\ell}(t)$ (5)

Their contribution can be integrated exactly: replace $\bigwedge_{\ell} (t)$ by $\bigwedge_{\ell} (t-\tau) = e^{-\frac{1}{\ell} t} \bigwedge_{\ell} (t)$ (6)

In practice, $e^{-iH_{\ell}^{1S}\tau}$ and $e^{iH_{\ell}^{b}\tau}$ are computed by using Krylov methods.

Build a Krylov space by applying $\mathcal{H}^{\text{IS}}_{\ell}$ multiple times to \mathcal{C}_{ℓ} , set up the tridiagonal representation $\mathcal{H}^{\text{IS}}_{\ell}$ in this basis, then compute the matrix exponential in this basis, and apply result to \mathcal{C}_{ℓ} . Likewise for $\mathcal{H}^{\text{IS}}_{\ell}$ and \mathcal{M}_{ℓ} .

To successively update entire chains, alternate between site- and bond-canonical form, propagating forward or backward in time with H_{p}^{15} or H_{p}^{b} , respectively:

1. Forward sweep, for
$$l=1,\ldots, l-1$$
 , starting from $B_1(t)B_2(t)\ldots B_\ell(t)$ (17)

$$C_{\ell}(t) \mathcal{B}_{\ell+1}(t)$$

$$\downarrow t + \tau \qquad A$$

$$\downarrow t \rightarrow t \rightarrow t$$

$$\downarrow t$$

until we reach last site, and MPS described by

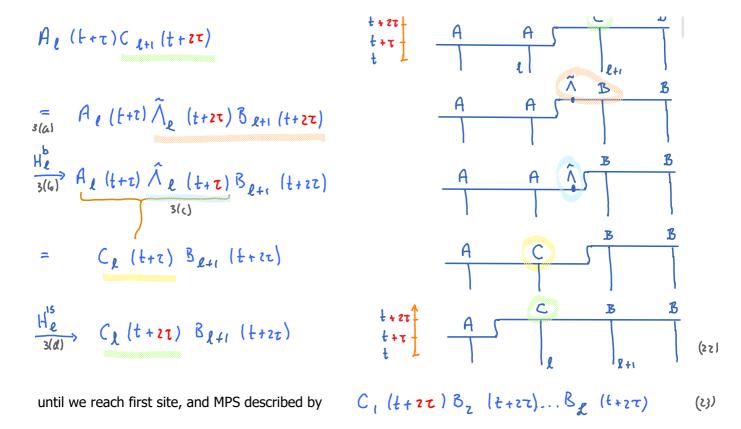
torf A A A 2. Turn around: (+)

$$\frac{H_{\mathcal{L}}^{15}}{2(a)} \left(\mathcal{L} \left(t + \tau \right) \right) \qquad \frac{A}{A} \qquad A \qquad A$$

$$\frac{H_{\mathcal{L}}^{IS}}{z(b)} \left(\left(t + 2\tau \right) \right) \quad \begin{array}{c} t + 2\tau \\ t + \tau \\ \end{array} \quad \begin{array}{c} C \\ \end{array}$$

3. Backward sweep, for $\ell = \mathcal{L} - 1$, ..., 1, starting from $A_1 (t+z) \dots A_{\ell-1} (t+z) C_{\ell} (t+zz)$ (2)

(20)



The scheme described above involves 'one-site updates'. This has the drawback (as in one-site DMRG), that it is not possible to dynamically exploring different symmetry sectors. To overcome this drawback, a 'two-site update' version of tangent space methods can be set up [Haegemann2016, App. C].

A systematic comparison of various MPS-based time evolution schemes has been performed in [Paeckel2019]. Conclusion: 2-site-update tangent space scheme is most accurate!

A scheme for doing 1-site TDVP while nevertheless expanding bonds, called 'controlled bond expansion (CBE), was proposed in [Li2022].

TS-II.2

The construction of tangent space V^{13} and its projector P^{13} can be generalized to n sites [Gleis2022a].

We focus on N = 7 (but general case is analogous). Define space of 2-site variations:

 \bigvee^{25} = span of all states \bigvee^{4} differing from \bigvee^{4} on precisely 2 neighboring sites

$$= span \left\{ \left| \vec{\Psi} \right\rangle = r \right\}$$

$$= \left\{ \left| \vec{\Psi} \right\rangle \right\} = r \left\{ \left| \vec{\Psi} \right\rangle \right\}$$

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formal definition: =
$$span \left\{ im \left(P_{\ell}^{zs} \right) \mid \ell \in [1, \ell-1] \right\}$$
 (2)

Recall:

Recall:
$$\frac{|\text{ocal}|}{|\text{local}|} \text{ 2s projector:} \qquad P_{\ell} = \text{(TS-I.4.9)}$$

Global 2s projector $\stackrel{\circ}{P}^{2s}$, such that $\stackrel{\vee}{V}^{2s} = i_{m} (\stackrel{\circ}{P}^{2s})$, can be found with a Gram-Schmidt scheme analogous to our construction of $\stackrel{\circ}{P}^{1s}$, see [Gleis2022a]:

compare (TS-I.5.22),
$$P^{2S} := \sum_{\ell=1}^{2S} P_{\ell}^{2S} + P_{\ell}^{2S} + \sum_{\ell=1}^{2S} P_{\ell}^{2S}$$

$$P^{2S} = \sum_{\ell=1}^{\ell'-1} \frac{1}{\ell} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'} \left\{ \begin{array}{c} \ell' \\ \ell+1 \end{array} \right\} + \frac{1}{\ell'}$$

All summands are mutually orthogonal, ensuring that $(p^{25})^2 = p^{25}$, and that $p^{25} = p^{25}$.

Alternative expression:

$$P^{25} = \sum_{\ell=1}^{l-1} P_{\ell}^{25} - \sum_{\ell=1}^{l-2} P_{\ell+1}^{15} = \sum_{\ell=1}^{l-1} \frac{1}{|\mathcal{L}|} \left(\frac{1}{|\mathcal{L}|} \right) \left(\frac{1}{|\mathcal{L}|$$

This projector is used for 2-site TDVP (see TS-II.3)

Orthogonal n-site projectors

, full Hilbert space of chain can be decomposed into mutually orthogonal subspaces: For any given MPS

$$V = V_1 \otimes \cdots \otimes V_\ell = \bigoplus_{N=0}^{\ell} V^{NL}$$
(8)

with
$$V^{ol} := V^{os} := Span \{12\}$$

'irreducible'
$$\bigvee^{N,\perp}$$
 is complement of $\bigvee^{(N-1)5}$ in $\bigvee^{N,5} = \bigvee^{(N-1)5} \bigoplus^{N} \bigvee^{N} \coprod^{(N-1)5} \bigoplus^{(N-1)5} \bigvee^{(N-1)5} \bigoplus^{(N-1)5} \bigoplus^{(N-1)5} \bigvee^{(N-1)5} \bigoplus^{(N-1)5} \bigoplus^{$

= span of states differing from $|\Psi\rangle$ on $|\Psi\rangle$ contiguous sites, not expressible through subsets of $|\Psi'\rangle$ sites

Correspondingly, identity can be decomposed as:

$$1_{V} = 1_{d}^{d} = \sum_{N=0}^{R} P^{NL}$$

$$= \sum_{N=0}^{N} P^{NL}$$

$$= \sum_{N=0}^{N} P^{NL}$$
orthogonality
$$= \sum_{N=0}^{N} P^{NL}$$
orthogonality

where P^{LN} is defined as the projector having V^{NL} as image: $Im(P^{NL}) = V^{NL}$ (12)

$$N \ge 1: \quad P^{NS} \left(\mathbb{1}_{V} - P^{(N-1)S} \right) = P^{NS} - P^{(N-1)S}$$

$$\text{since im} \left(\mathbb{V}^{(N-1)S} \right) \subset \text{im} \left(\mathbb{V}^{NS} \right)$$

Consider n=1:

$$= \sum_{\ell=1}^{\ell'} \sum$$

 $\operatorname{im}(P^{(n-1)S}) \subset \operatorname{im}(P^{NS})$

(2)

choose l' = L

$$= \sum_{\ell=1}^{l} P_{\ell,\ell+1}^{bk}$$
 projects onto all 1-site variations orthogonal to $|\Psi\rangle$

- **

Consider n=2:

$$P^{2,1} = P^{2,5} - P^{1,5} = \begin{pmatrix} \lambda^{-1} & P^{2,5} & -\lambda^{-1} & P^{1,5} \\ \lambda^{-1} & \lambda^{-1} & \lambda^{-1} \end{pmatrix} - \begin{pmatrix} \lambda^{-1} & P^{1,5} & -\lambda^{-1} & P^{1,5} \\ \lambda^{-1} & \lambda^{-1} & \lambda^{-1} & \lambda^{-1} \end{pmatrix}$$

$$= \begin{pmatrix} \lambda^{-1} & P^{1,5} & -\lambda^{-1} & P^{1,5} \\ \lambda^{-1} & \lambda^{-1} & \lambda^{-1} & \lambda^{-1} \end{pmatrix}$$

$$= \begin{pmatrix} \lambda^{-1} & \lambda^{-1} & \lambda^{-1} & \lambda^{-1} \\ \lambda^{-1} & \lambda^{-1} & \lambda^{-1} & \lambda^{-1} \end{pmatrix}$$

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$$= \begin{pmatrix} \lambda^{-1} & \lambda^{-1} & \lambda^{-1} & \lambda^{-1} & \lambda^{-1} & \lambda^{-1} \\ \lambda^{-1} & \lambda^{-1} & \lambda^{-1} & \lambda^{-1} \end{pmatrix}$$

$$= \begin{pmatrix} \lambda^{-1}$$

$$= \begin{array}{c} \mathcal{L}^{-1} \\ = \\ \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array} \qquad = \begin{array}{c} \mathcal{L}^{-1} \\ \ell = 1 \end{array}$$

[Haegeman2016, Sec. V & App. C]

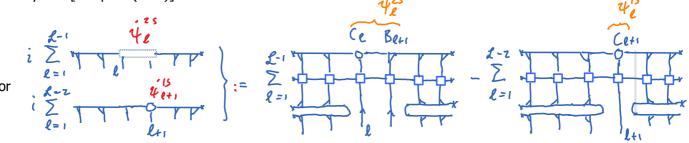
2-site tangent space methods are analogous to 1-site methods, but use a 2-site projector. There is a conceptual difference, though: the main reason for using 2-site schemes is that they allow sectors with new quantum numbers to be introduced if the action of H requires this. However, states with different ranges of quantum numbers live in different manifolds, hence this procedure 'cannot easily be captured in a smooth evolution described using a differential equation. However, like most numerical integration schemes, the aforementioned algorithm is intrinsically discrete by choosing a time step, and it poses no problem to formulate an analogous two-site algorithm'. [Haegeman2016, Sec. V]. In other words: the tangent space approach is conceptually not as clean for the 2-site as for the 1-site scheme.

Schrödinger equation, projected onto 2-site tangent space, now takes the form

$$i \frac{d}{at} | \psi[m(t)] \rangle = \hat{\rho}^{2s} \hat{H} | \psi[m(t)] \rangle$$

$$\hat{P}^{ZS} = \sum_{\ell=1}^{2^{-1}} \frac{1}{\ell+1} \left| \frac{1}{\ell+1} - \sum_{\ell=2}^{2^{-1}} \frac{1}{\ell+1} \right|$$

This yields [compare (1.11)]:



Right side is sum of terms, each specifying an update of one ψ_{ℓ}^{ts} or ψ_{ℓ}^{ts} on the left. Eq. (4) can be integrated one site at a time, by defining the updates through the following local Schrödinger equations:

$$\frac{i \dot{\psi}_{\ell}^{2S}}{!} :=
\begin{array}{c}
\psi_{\ell}^{1S} \\
\psi_{\ell}^{1S}
\end{array}$$

$$\frac{i \dot{\psi}_{\ell+1}^{1S}}{!} :=
\begin{array}{c}
\psi_{\ell+1}^{1S} \\
\psi_{\ell+1}^{1S}
\end{array}$$

Right side is sum of terms, each linear in a factor appearing on the left. Can be integrated one site at a time:

In 2-site-canonical form, site
$$\ell$$
 involves two terms linear in Ψ_{ℓ}^{zs} : $i \Psi_{\ell}^{zs}(t) = H_{\ell}^{zs} \Psi_{\ell}^{zs}(t)$ (1)

Their contribution can be integrated exactly: replace
$$\psi_{\ell}^{2s}(t)$$
 by $\psi_{\ell}^{2s}(t+\tau) = e^{-i H_{\ell}^{2s} \tau} \psi_{\ell}^{2s}(t)$ forward time step

In 1-site-canonical form, site
$$\ell$$
+1 involves two terms linear in $\Psi_{\ell+1}^{15}$: $i \Psi_{\ell+1}^{15}(t) = - H_{\ell+1}^{15} \Psi_{\ell+1}^{15}(t)$ (12)

Their contribution can be integrated exactly: replace
$$\psi_{\ell+1}^{(s)}(t)$$
 by $\psi_{\ell+1}^{(s)}(t-\tau) = e^{iH_{\ell+1}^{(s)}\tau} \psi_{\ell+1}^{(s)}(t)$ (3)

Their contribution can be integrated exactly: replace
$$\psi_{\ell+1}^{(s)}(t)$$
 by $\psi_{\ell+1}^{(s)}(t-\tau) = e^{iH_{\ell+1}^{(s)}\tau} \psi_{\ell+1}^{(s)}(t)$ backward(!) time step

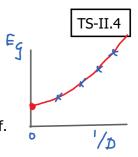
To successively update entire chains, alternate between 2-site- and 1-site-canonical form, propagating forward or backward in time with H_{ℓ}^{25} or H_{ℓ}^{15} , respectively (analogously to 1-site scheme).

A systematic comparison of various MPS-based time evolution schemes has been performed in [Paeckel2019]. Conclusion: 2-site-update tangent space scheme is most accurate!

4. Energy variance

[Hubiq2018]

When doing MPS computations involving SVD truncations of virtual bonds, the results should be computed for several values of the bond dimension, 70, to check convergence as $\triangleright \rightarrow \infty$. Often it is also necessary to extrapolate the results to $\mathbb{D} = \emptyset$, e.g. by plotting results versus $1/\mathbb{D}$ or some power thereof.

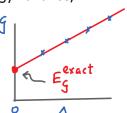


However, for some computational schemes, it is not a priori clear how the observable of interest scales with \mathcal{D} , nor how it should be extrapolated to $\mathcal{D} = \infty$. An example is ground state energy when computed using 1-site DMRG with subspace expansion [Hubig2015], because it does not rely on SVD truncation of bonds.

Thus, it is of interest to have a reliable error measure without requiring costly 2-site DMRG. A convenient scheme was proposed in [Hubig2018], based on a smart way to approximate the full energy variance,

$$\Delta_{E} := \| (H - E) \psi \|^{2} = \langle \psi | (H - E)^{2} | \psi \rangle \quad (= \text{zero for an exact eigenstate}) \quad (1)$$

$$= \langle \psi | H^{2} | \psi \rangle - E^{2}, \quad \text{with} \quad E = \langle \psi | H^{2} | \psi \rangle \quad (2)$$



Then extrapolations can be done by computing quantity of interested for several \mathbb{D} , but plotting the results via $\ \triangle_{\mathsf{E}}\$, and extrapolating to $\ \triangle_{\mathsf{E}}\ o$ $\ o$

If quantity of interest is energy, then extrapolation is linear, $\varepsilon_{g}(\Delta_{\varepsilon}) = \varepsilon_{c}^{exect} + \alpha \cdot \Delta_{\varepsilon}$

$$E_g(\Delta_E) = E_g^{exect} + a \cdot \Delta_E$$
 (3)

Computing $\langle \psi | \hat{\mu}^{\dagger} | \psi \rangle$ directly is costly for large systems with long-ranged interactions, such as 2D systems treated by DMRG snakes. Also, computing \triangle_{ϵ} as the difference between two potentially large numbers is prone to inaccuracies. [Hubig2018] found a computation scheme in which the subtraction of such large numbers is avoided a priori.



Key idea: use projectors → onto mutually orthogonal, irreducible spaces

Recall (2.11):
$$1_{V} = 1_{d}^{\otimes \ell} = \sum_{N=0}^{\ell} P^{N\perp} \qquad P^{N} = \sum_{N=0}^{\ell} P^{N\perp} \qquad \text{orthogonality}$$
 (5)

with
$$P^{\circ \perp} = |\Psi\rangle\langle\Psi|$$
 (6)

Insert completeness into definition of variance:
$$\Delta_{\epsilon} = \langle \psi | (\hat{\mu} - \epsilon) \sum_{N=0}^{\ell} P^{NL}(\hat{\mu} - \epsilon) | 14 \rangle = : \sum_{N=0}^{\ell} \Delta_{\epsilon}^{NL}$$
 (8)

Now two crucial simplifications occur:

$$\Delta_{E}^{OL} = \langle \psi | (\hat{H} - E) | \psi \rangle \langle \psi | (\hat{H} - E) \rangle | \psi \rangle = (E - E) \langle E - E \rangle = 0$$
[4]

Ignorest contribution to variance cancels by construction]

$$\Delta_{E}^{OL} \stackrel{(5)}{=} \langle \psi | (\hat{H} - E) | \psi \rangle \langle \psi | (\hat{H} - E) \rangle | \psi \rangle = (E - E) \langle E - E \rangle = 0$$
(q)
largest contribution to variance cancels by construction!

 $\nabla_{\mathsf{nT}}^{\mathsf{E}} = \langle \mathsf{A} | (\mathring{\mathsf{H}} - \mathsf{E}) \rangle_{\mathsf{nT}} (\mathring{\mathsf{H}} - \mathsf{E}) | \mathsf{A} \rangle = \langle \mathsf{A} | \mathring{\mathsf{H}} \rangle_{\mathsf{nT}} \mathring{\mathsf{H}} | \mathsf{A} \rangle^{\mathsf{L}} \text{ since } \mathcal{b}_{\mathsf{(N>0)}} \mathsf{T} | \mathsf{A} \rangle = 0 \quad (10)$

$$= \|p^{n\perp} \hat{p}\| \Psi\|^2$$

In practice, approximate $\triangle_{\mathcal{E}}$ by the first two nonzero terms:

$$\Delta_{\varepsilon} \simeq \Delta_{\varepsilon}^{\varepsilon} = \Delta_{I}^{\varepsilon} + \Delta_{\zeta}^{\varepsilon} = \left\langle \chi_{I} | \hat{H} P_{I}^{\varepsilon} \hat{H} | \chi \right\rangle + \left\langle \chi_{I} | \hat{H} P_{\zeta}^{\varepsilon} \hat{H} | \chi \right\rangle \quad (15)$$

(11) is exact if longest-range terms in \hat{H} are nearest-neighbor, because then $p(N \geqslant 3) \perp \hat{H} \mid \psi \rangle = 0$ (13) Explicit computations:

$$N = 1: \text{ Recall } P^{1} = \sum_{\ell=1}^{(2.16)} \frac{\ell}{\ell} = \sum_{\ell=1}^{(2.16)} P^{1}_{\ell} = \sum_{\ell=1}^{(2.16)} P$$

$$\Delta_{E}^{1L} = \langle \psi | \hat{H} \hat{P}^{1L} \hat{H} | \psi \rangle = ||P^{1L} H \psi||^{2} = \sum_{l=1}^{2} ||P_{l,l+1}^{DK} H \psi||^{2}$$
 (15)

$$=\sum_{\ell=1}^{2}\sum_{\ell=1}^$$

We would like to avoid computing explicitly, because of its large image dimension.

So rewrite, using isometry condition for discarded sector:

and completeness of kept together with discarded isometries: = -

$$= \sum_{\ell=1}^{2} \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right| = \sum_{\ell=1}^{2} \left| \begin{array}{c} 1 \\ 1 \end{array} \right$$

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$$N = z$$
: Recall $P^{z} = \sum_{\ell=1}^{(2.17)} \frac{\ell}{\ell} = \sum_{\ell=1}^{(2.17)} P_{\ell,\ell+1}^{z} = \sum_{\ell=1}^{(2.1$

$$\Delta_{E}^{z} = \langle 4 | \hat{H} P^{z} \hat{H} | 4 \rangle = ||P^{z} H \Psi||^{2} = \sum_{l=1}^{\infty} ||P^{l}_{l,l+1} H \Psi||^{2}$$
 (21)

$$= \begin{array}{c|c} \mathcal{L}^{-1} & & \\ & & \\ & & \\ & & \\ & & \\ \end{array}$$

$$= \begin{array}{c|c} \mathcal{L}^{-1} & & \\ & & \\ & & \\ & & \\ \end{array}$$

$$= \begin{array}{c|c} \mathcal{L}^{-1} & & \\ & & \\ & & \\ & & \\ \end{array}$$

$$(72)$$

again use
$$\Rightarrow = \Rightarrow | - \Rightarrow |$$
 (23)