DMRG-II.1

1. Original formulation of DMRG

Goal: finding ground state of infinite chain

Infinite-size DMRG (iDMRG)

Diagonalize small system (e.g. 2 sites), write ground state in the form

$$|\psi\rangle = \sum_{\alpha\beta} |\beta\rangle_R |\alpha\rangle_L 2^{\alpha\beta}$$

Now add two sites between blocks L and R, and seek new ground state of $H_{L^{\bullet}R}$ of the form

$$|\psi\rangle = |\beta\rangle_{R} |\epsilon_{R}\rangle |\epsilon_{L}\rangle |\kappa\rangle_{L} \psi^{\kappa} |\psi\rangle$$

by minimizing (Lanzcos)
$$\frac{(\psi|\hat{H}_{L}, \kappa|\psi)}{(\psi|\psi)}$$



Write ground state in the form

(2)

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Then iterate,: add two more sites, etc.



Remark: we established early on (see MPS-II.2) that the eigenvalues ρ_{c} of reduced density matrix of L, are obtained by SVD of ψ^{ab}



So retaining the β_L -eigenstates with largest ρ_c is equivalent to just doing SVD-truncation on ψ^{ab}

Modern formulation

Start with MPS in bond-canonical form:

Add two central sites and find ground state (Lanczos):

Do SVD to split chain into two larger blocks, and truncate: $\Psi = \mu S V^{+} = \hat{A} \hat{\chi} \hat{\delta}$



Iterate: make chain longer and longer, until ground state energy per site converges.

'iDMRG state prediction' [McCulloch2008], [Schollwöck2011, Sec. 10.1]

To speed up Lanczos search for ground state, construct initial guess for ψ from previous data:



This leads to 'dramatic speedup' of iDMRG.

Finite-size DMRG

Grow chain to some length N using infinite-size DMRG algorithm.

Then reduce L, enlarge R, optimize ψ :

Diagonalize ρ_{L} , truncate.



This is conceptually <u>identical</u> to variational optimization with two-site update.

Single-site DMRG is also possible \iff variational single-site update.

Problem: when exploiting quantum numbers, 1-site DMRG performs poorly, because it does not explore subspaces with different quantum numbers. An early remedy for this is 2-site DMRG, but that is computationally much more expensive than 1-site DMRG. Subsequent suggestions for 1-site DMRG with symmetries are 'density matrix perturbation' [White2005], the 'center matrix wave function formalism [McCulloch2007], and, most recently, 'subspace expansion' [Hubig2015], which performs best.

fixed normalization, 1 site at a time:



Minimize energy with constraint of fixed normalization, 1 site at a time: $\frac{\partial}{\partial B_{11}^{\dagger}} \left[\langle 4 | \hat{H} | 4 \rangle - \lambda \langle 4 | 4 \rangle \right] = 0$ (2)



$$\alpha = (\nu, \delta, \beta)$$

Solve for 'eigenvector' with lowest eigenvalue, say $M_{[\ell]}$, then do SVD on it to move to next site:

$$D \xrightarrow{\widetilde{M}[e]}_{d} \xrightarrow{B[e+i]}_{d} \xrightarrow{SVD}_{d} \xrightarrow{U} \underbrace{(U)(S V^{\dagger} B[e+i])}_{d} \xrightarrow{D} \xrightarrow{D} \underbrace{\widetilde{A}[e1 M_{[e+i]}]}_{d} \xrightarrow{(S)} \xrightarrow{(S)}_{d} \xrightarrow{D} \xrightarrow{\widetilde{A}[e1 M_{[e+i]}]}_{d} \xrightarrow{(S)} \xrightarrow{(S)}_{d} \xrightarrow{D} \xrightarrow{\widetilde{A}[e1 M_{[e+i]}]}_{d} \xrightarrow{(S)}_{d} \xrightarrow{(S)}_{d}$$

Important: dimensions of $\widetilde{M}_{[\ell]}$ are <u>fixed</u>, hence truncation is neither needed nor possible!

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Problem of single-site optimization: it is constrained to a variational space defined by outgoing state spaces β_{ℓ_1} α_{ℓ_2} . If the ranges of quantum numbers β_{κ} and β_{β} for these spaces are too small to accurately represent the ground state, single-site DMRG has no way to enlarge them.

Two-site optimization does not have this problem: the action of H on two sites enlarges bond dimension in between, adding the full range of quantum numbers needed on that bond. If a certain quantum number was missing on that bond before the action of H, but appears afterwards with non-negligible weight, it will survive after SVD and truncation. Hence: two-site optimization can add missing quantum numbers, if needed.

But this comes at a cost: effective two-site Hamiltonian has dimension $D^2 d^2 \times D^2 d^2$. (9) By contrast, effective one-site Hamiltonian has dimension $D^2 d \times D^2 d$.

(13)

Example: 6-site chain of S=1/2 spins, exploiting Sz quantum numbers: $\oint e^{\frac{1}{2}} e^{\frac{1}{2}} \frac{1}{2} \frac$

$$\begin{split} & (\psi) = \frac{A}{k} - \frac{A}$$

In general, the product $M_{(3)} \mathcal{B}_{[4]}$ has block matrix product structure, where sum over bond index $\beta = (Q_3, i_3)$ matches blocks with same Q_3 :



 $D_2 \cdot d_3 \times D_3 = 4 \cdot 2 \times 8$ $D_3 \times d_4 \cdot D_4 = 8 \times 2 \cdot 4$ Bond dimensions:

But in practice, we use truncated state spaces. Suppose $\mathcal{D}_{MA_{X}} = 4$,

with $D_z = 4$, $d_3 = 2$, $D_3 = 4$ and that random initial state does not contain $Q_3 = \pm 3$ Then the corresponding two blocks are missing in both $M_{[3]}$ and $\mathcal{B}_{[4]}$:

Truncated:

1-site optimization of truncated $M_{[3]}$ will never find a good ground state if the latter has non-negligible contributions from missing blocks.

2-site optimization reinstates missing block: E.g.:



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for each fixed μ :

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$$= \left(\begin{array}{c} L_{\left[\ell - i \right]} \\ W_{\left[\ell \right]} \\ M_{\left[\ell \right]} \end{array} \right)^{\alpha} \left(\begin{array}{c} z \\ z \\ z \end{array} \right) = \left(\begin{array}{c} B_{\left[\ell + i \right]} \\ W_{\left[\ell + i \right]} \\ R_{\left[\ell + i \right]} \\ \beta \\ \beta \\ \mu \end{array} \right)^{\overline{\beta}} \left(\begin{array}{c} z \\ z \\ z \end{array} \right)$$
dimensions: $\left(\begin{array}{c} D \\ i \\ d \end{array} \right)^{\alpha} \left(\begin{array}{c} d \\ w \\ \overline{D} \end{array} \right) \left(\begin{array}{c} z \\ z \\ z \end{array} \right)$
dimensions: $\left(\begin{array}{c} D \\ i \\ d \end{array} \right)^{\alpha} \left(\begin{array}{c} d \\ w \\ \overline{D} \end{array} \right) \left(\begin{array}{c} z \\ z \\ z \end{array} \right)$

Now we add the two MPS in (16), by yet again expanding virtual state space and block-multiplying their tensors:

$$\begin{split} |\tilde{\psi}\rangle &= |\psi_{\sigma}\rangle + c \langle \psi_{\alpha}\rangle = |\alpha, \sigma, \bar{\sigma}, \chi\rangle \left[(M_{[\ell]})^{\kappa \sigma \beta} (B_{[\ell+1]})_{\beta} \bar{\sigma}^{\gamma} + c \left[(L_{[\ell]})^{\alpha \sigma, \mu \beta} (\tilde{R}_{[\ell+1]})_{\beta \mu} \bar{\sigma}^{\gamma} \right] (ze) \\ \left[M_{(e]} - c \overline{L}_{[\ell]} \right]^{\kappa \sigma} \bar{\beta} \cdot \left[B_{[\ell+1]} \right]_{\beta} \bar{\sigma}^{\gamma} =: (\mathcal{M} \cdot \mathcal{B})^{\alpha \sigma \bar{\sigma} \gamma} (25) \\ (16) M_{(e]} - c \overline{L}_{[\ell]} \int_{\beta} \bar{\beta} (D_{\ell}) \cdot d(D_{\ell}) \times (D_{\ell}) (26) \\ dim(\tilde{\beta}) &= dim(\beta) + dim(\tilde{\beta}) \cdot dim(\mu) (26) \end{split}$$



. . . <u>.</u>

To deal with large-dimension bond β , use SVD on $\mathcal{M} = \mathcal{U} S V^{\dagger}$ and <u>truncate</u> back to (!!)

$$\frac{\mathcal{M}}{\mathcal{D}} = \frac{\mathcal{B}}{\mathcal{D}} = \frac{\mathcal{M}}{\mathcal{D}} = \frac{\mathcal{M}}{\mathcal{M}} = \frac{\mathcal{M}}{\mathcal{M}$$

Left-normalized $A_{[\ell]}$ is final result of optimization of site ℓ , and $M_{[\ell]}$ initializes optimization of site $\ell + \ell$.

Crucial point: as mentioned earlier, β may include larger range of quantum numbers than β . If these contribute significant weight to ψ_{α} , then they will have significant weight in SVD of (27), hence will survive the truncation in (27), and thus be present in \tilde{A}_{β} .

In this way, subspace expansion succeeds in 'reinstating missing quantum numbers' (without ever needing to do Lanczos with huge effective 2-site Hamiltonian of 2-site DMRG!).

6-site example: if $Q_3 = \pm 3$ carries significant weight, then SVD, followed by trunction with dim(β) = 4, will reinstate those blocks:



Remark: [Hubig2015] choose $\overline{\mathcal{R}}_{[\ell+1]} = \mathfrak{o}$ [instead of (22b)], arguing that it only affects $\mathcal{M}_{[\ell+1]}$, which will be updated anyway when optimizing site $\ell+1$ via Lanczos ground state solver in next iteration. But using (21b) speeds up convergence for that Lanczos solver [Weichselbaum (private communication)].

<u>Summary of subspace expansion algorithm</u> 'strictly single-site DMRG (DMRG3S)'

Left-to-right sweep:

1. In site-canonical basis for site
$$\ell$$
, use 1-site DMRG to find optimized $M_{[\ell]}$
with $H_{[\ell]} M_{[\ell]} = E_{[\ell]} M_{[\ell]}$ (c1)
new current energy estimate
2. Expand subspace, $M_{[\ell]} \beta_{[\ell rr]} \rightarrow M_{[\ell]} \cdot \overline{\beta}_{[\ell rr]} = \left(M_{[\ell]} - \overline{C}_{[\ell]}\right) \begin{pmatrix} \beta_{[\ell rr]} \\ \overline{K}_{[\ell rr]} \end{pmatrix} \approx \widehat{A}_{[\ell]} M_{[\ell rr]} (3n)$
3. Obtain 'post-expansion' energy estimate $\widetilde{E} = \underbrace{\prod_{l \in I} M_{[\ell r]} \\ M_{[\ell r]} \\ \overline{n}_{l \in I} \\$

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3. Obtain 'post-expansion' energy estimate

Choice of mixing coefficient **c** :

usually < 0

SET.



Behavior of energy changes depends on C

can be either < 0 or > 0.

energy change from truncation,

If c is too small, one can get stuck in local minimum, i.e. $\Delta E_0 \simeq c$

If ζ is too large, truncation destroys changes made by previous local optimization, and $0 \leq \Delta E_{\tau} \approx |\Delta E_{0}|$ (34) Then reduce *c* for next iteration on next site.

If ΔE_{T} is much smaller than $|\Delta E_{o}|$ or even negative, then one can increase c for next iteration.

In practice, start with c = 1 and aim for $\Delta \epsilon_{\tau} \simeq 0.3 | \Delta \epsilon_{o} |$ (for scheme of [Hubig2015], with $\overline{R} = 0$) and change C by multiplying with some factor greater/smaller than 1 (rather than subtracting/adding). After several DMRG sweeps back and forth through the chain, the 'correct' quantum numbers will typically have been found. So, one can eventually take $\subset \rightarrow \circ$, returning to pure 1-site DMRG.

Remark: there may be better ways to choose c : build the 2-dimensional Krylov space $K = \{1, 4, 5, 1, 4, 5\}$ $|\hat{\psi}\rangle = |\psi_0\rangle + c|\psi_0\rangle$ is the ground state of in this space. and choose *c* such that (35)This corresponds to the first step of a Lanczos iteration, cf. DMRG-I.2, Eqs. (6-11).

 $b_1 | \psi_1 \rangle := | \psi_1 \rangle := \hat{p} \hat{H} | \psi_0 \rangle - | \psi_0 \rangle \langle \psi_0 | \hat{p} \hat{H} | \psi_0 \rangle$ (36) Orthogonalize: = 14a) - ao 140) (37)

Computing $|\hat{\psi}\rangle$ involves addition of two MPS, can be done as in (24,25) [no need for SVD + truncation here], to arrive at an MPS with same bond dimensions as $|\psi_0\rangle$.

$$b_1 = \langle \widetilde{\psi}_1 | \widetilde{\psi}_1 \rangle = \langle \psi_1 | H | \psi_0 \rangle, \quad a_1 = \langle \psi_1 | H | \psi_1 \rangle \quad (38)$$

 $H_{\kappa} = \begin{pmatrix} q_{\circ} & b_{\circ} \\ b_{\circ} & q_{\circ} \end{pmatrix}$, its lower-energy eigenvector yields optimal c. Then diagonalize (39) to compute a new $|\hat{\psi}\rangle$, and from that compute $\mathcal{M} - \mathcal{B} = \widehat{A}_{[\ell]} \mathcal{M}_{[\ell+1]}$ Then use that as described in (24,25,27). Explore whether this yields better convergence behavior!

How can one estimate / control error during extrapolations with increasing bond dimension? See [Hubig2018].