<\varphi | \psi > 1. Overlaps, matrix elements

We first consider general quantum states, then matrix product states (MPSs):

 $|\psi\rangle = |\epsilon_{N}\rangle \dots |\epsilon_{2}\rangle|\epsilon_{1}\rangle C^{\epsilon_{1},\dots,\epsilon_{N}} := |\epsilon\rangle C^{\epsilon_{1}}$ General ket: (E 44N)

(1)

summation over repeated indices implied

General bra: $\langle \psi | = C_{6_1, \dots, 6_N} \rangle \langle 6_1 | \langle 6_2 | \dots | \langle 6_N | \rangle = C_{6_R} \langle \delta | \sum_{i=1}^{n} C_{6_R} \rangle \langle \delta | \sum_{i=1}^{n} C_{6_R} \langle \delta | \sum_{i=1}^{n} C_{6_R} \rangle \langle \delta | \sum_{i=1}^{n} C_{6_R} \langle \delta | \sum_{i=1}^{n} C_{6_R} \rangle \langle \delta | \sum_{i=1}^{n} C_{6_R} \langle \delta | \sum_{i=1}^{n} C_{6_R} \rangle \langle \delta | \sum_{i=1}^{n} C_{6_R} \langle \delta | \sum_{i=1}^{n} C_{6_R} \rangle \langle \delta | \sum_{i=1}^{n} C_{6_R} \langle \delta |$

(2)

(ψ | ψ) = (6,1,...,6μ (6,1) (6,1)... (6,1) (6,1) ... (6,2) | 6,2) | 6,1,..., 6μ (3a) 86, 80263 80% = 8 0 0

 $= \tilde{C}_{\vec{\sigma}_{\sigma}}^{\dagger} C^{\vec{\sigma}}$

(3b)

Recipe for overlaps: contract all physical legs of bra and ket.

General operator:

0 = 160 00 = (6)

(4)

Matrix

Matrix elements: $\langle \hat{\psi} | \hat{\sigma} | \psi \rangle = C \vec{\sigma} \langle \vec{\sigma} | \vec{\sigma} \rangle C \vec{\sigma} \langle \vec{\sigma} | \vec{\sigma} \rangle C \vec{\sigma}$

= (+ 0 + 0

(5a) (5b)

Recipe for matrix elements: contract all physical legs of bra and ket with operator.

Now consider matrix product states:

Ket: 14>=19> 4 [1] x 4 [5] & 4 [3] 4... 4 [1]

ALI) A(a) A(a) Acura

Recipe for ket formula: as chain grows, attach new matrices A^{6} on the right (in <u>same</u> order as vertices in diagram), resulting in a matrix product of H^{6} matrices.

index-reading order

Square brackets indicate that each site has a different A^{σ} matrix. We will often omit them and A « G = A(l) B. since the ℓ on σ_{ℓ} uniquely identifies the site. use the shorthand,

Add dummy sites at left and right, so that first and last A's have two virtual indices, just like other A's .

Bra:
$$\langle \psi | = \overrightarrow{A_{[i]}^{1\varsigma_{i}}} \stackrel{\wedge}{A_{[i]}^{\alpha}} \stackrel{\wedge}{A_{[i]}^{\beta\varsigma_{3}}} \stackrel{\wedge}{\wedge} \stackrel{\wedge}{A_{[i]}^{\beta\varsigma_{1}}} \stackrel{\wedge}{\wedge} \stackrel{\wedge}{\delta} \stackrel{\wedge}{\wedge} \stackrel{$$

We expressed all matrices via their Hermitian conjugates by transposing indices and inverting arrows. To recover a matrix product structure, we ordered the Hermitian conjugate matrices to appear in the opposite order as the vertices in the diagram.

Recipe for bra formula: as chain grows, attach new matrices $\mathbf{A}_{\mathbf{S}}^{\mathbf{t}}$ on the left, (in opposite order as vertices in diagram), resulting in a matrix product of A_{\bullet}^{\dagger} matrices.

Overlap:
$$\langle \psi | \psi \rangle \stackrel{(3b)}{=} \stackrel$$

Recipe: contract all physical indices with each other, and all virtual indices of neighboring tensors.

Matrix elements:
$$\langle \hat{\mathcal{A}} \mid \hat{0} \mid \psi \rangle =$$

$$A_{[i]} = A_{[i]}^{[i]} A_{[i]}^{[i]$$

Exercise: derive this result algebraically from (7a), (8a)!

If we would perform the matrix multiplication first, for fixed \vec{r} , and then sum over \vec{c} , we would get d^{N} terms, each of which is a product of 2N matrices. Exponentially costly!



But calculation becomes tractable if we rearrange summations:

$$\langle \psi | \psi \rangle = C_{[i]} \qquad C_$$

$$= \widetilde{A}^{\dagger}_{\delta_{N,M}} \dots \widetilde{A}^{\dagger}_{\delta_{2M}} A^{\dagger}_{\delta_{1}} \dots A^{\dagger}_{[i]} A^{\dagger}_{\delta_{1}} \dots A^{\dagger}_{[i]} A^{\dagger}_{\delta_{2M}} \dots A^{\dagger}_{[N]}$$

$$:= C^{\dagger}_{[i]} A$$

Diagrammatic depiction: 'closing zipper' from left to right.

The set of two-leg tensors $C_{[\ell]}$ can be computed iteratively:

Initialization:
$$C_{\{0\}} = C_{\{0\}} = C_{\{0\}}$$

Iteration step:
$$C_{\{\ell\}} = C_{\{\ell-1\}} + C_{\{\ell\}} + C_{\{\ell-1\}} + C_{\{$$

Cost estimate (if all A's are $\mathcal{D}_{\chi} \mathcal{D}$):

Page 3



Total cost: $\sim D^3 d \cdot N$

Remark: a similar iteration scheme can be used to 'close zipper from right to left':

$$D_{[N+1]} = C_{[N-1]} D_{[N]} = C_{[N]} D_{[N]} = C_{[N]} D_{[N]} D_{[N]} = C_{[N]} D_{[N]} D_{[N]}$$

Normalization $\langle \psi | \psi \rangle = ?$ Use above scheme, with $\hat{A} = A$

Computation of expectation and matrix elements is simpler if the MPS is built from tensors relating orthonormal spaces. Such tensors are called 'left-normalized' or 'right-normalized.

Left-normalization

A 3-leg tensor $A^{\alpha \delta}$ is called 'left-normalized' if it satisfies

$$A^{\dagger}A = 1$$
 Explicitly: $(A^{\dagger}A)^{\beta'}_{\beta} = A^{\dagger}\beta'_{\delta\alpha}A^{\alpha\delta}_{\beta} = 1^{\beta'}_{\beta}$ (1)

Graphical notation for left-normalization: draw 'left-pointing diagonals' at vertices

When all A's are left-normalized, closing the zipper left-to-right is easy, since all reduce to identity matrices:

$$C_{[o]} = \begin{cases} c_{[i]_{\alpha}} = c_{[a]_{\alpha}} \\ c_{[i]_{\alpha}} = c_{[a]_{\alpha}} \end{cases}$$

$$C_{[o]} = \begin{cases} c_{[a]_{\alpha}} = c_{[a]_{\alpha}} \\ c_{[a]_{\alpha}} = c_{[a]_{\alpha}} \end{cases}$$

$$C_{[o]} = \begin{cases} c_{[a]_{\alpha}} = c_{[a]_{\alpha}} \\ c_{[a]_{\alpha}} = c_{[a]_{\alpha}} \end{cases}$$

Hence:

$$\langle \psi(\psi) \rangle = \left(\begin{array}{c} \psi(\psi) \\ \psi(\psi) \end{array} \right) = \left(\begin{array}{c} \psi(\psi) \\ \psi(\psi) \\ \psi(\psi) \end{array} \right) = \left(\begin{array}{c} \psi(\psi) \\ \psi(\psi) \\ \psi(\psi) \end{array} \right) = \left(\begin{array}{c} \psi(\psi) \\ \psi(\psi) \\ \psi(\psi) \\ \psi(\psi) \end{array} \right) = \left(\begin{array}{c} \psi(\psi) \\ \psi$$

When all matrices of a MPS are left-normalized, the matrices for site 1 to any site $\ell = 1, ..., N$ define an orthonormal state space:

$$|\lambda\rangle = |\delta_{\ell}\rangle |A^{\delta_{1}} A^{\delta_{2}} ... A^{\delta_{\ell}}|^{1}$$

$$|\lambda\rangle = |\delta_{\ell}\rangle |A^{\delta_{1}} A^{\delta_{2}} ... A^{\delta_{\ell}}|^{1}$$

$$|\lambda\rangle = |\lambda\rangle |A^{\delta_{1}} A$$

Right-normalization

So far we have viewed an MPS as being built up from left to right, hence used right-pointing arrows

on ket diagram. Sometimes it is useful to build it up from right to left, running left-pointing arrows.

Building blocks:

$$\langle \beta | = \mathcal{B}_{16N}^{\dagger} \mathcal{B}_{\alpha 6N-1}^{\dagger} \langle 6N-1 | \langle 6N | \mathcal{B}_{\alpha 6N-1}^{\dagger} \rangle$$
 (10)

Iterating this, we obtain kets and bras of the form

$$|\psi\rangle = |\epsilon_{N}\rangle|\epsilon_{N-1}\rangle...|\epsilon_{1}\rangle B_{1}^{\epsilon_{1}\lambda}...B_{n}^{\epsilon_{N-1}}B_{n}^{\alpha}B_{n}^{\alpha}$$

$$= |\epsilon_{N}\rangle|\epsilon_{N-1}\rangle...|\epsilon_{1}\rangle B_{1}^{\epsilon_{1}\lambda}...B_{n}^{\epsilon_{N-1}}B_{n}^{\alpha}$$

$$= |\epsilon_{N}\rangle|\epsilon_{N-1}\rangle...|\epsilon_{N-1}\rangle...B_{n}^{\epsilon_{N-1}}B_{n}^{\alpha}$$

$$= |\epsilon_{N}\rangle|\epsilon_{N-1}\rangle...B_{n}^{\epsilon_{N-1}}B_{n}^{\alpha}$$

$$= |\epsilon_{N}\rangle|\epsilon_{N-1}\rangle...B_{n}^{\epsilon_{N-1}}B_{n}^{\alpha}$$

$$= |\epsilon_{N}\rangle|\epsilon_{N-1}\rangle...B_{n}^{\epsilon_{N-1}}B_{n}^{\alpha}$$

$$= |\epsilon_{N-1}\rangle|\epsilon_{N-1}\rangle...B_{n}^{\epsilon_{N-1}}B_{n}^{\alpha}$$

$$\langle 4 | = B_{16N}^{\dagger} B_{\alpha 6_{N-1}}^{\dagger} B_{\alpha 6_{N-1}}^{\dagger} B_{\alpha 6_{N-1}}^{\dagger} \langle 6_{1} | \dots \langle 6_{N-1} | \langle 6_{N} | \dots \rangle \langle 6_{N-1} | \langle 6_{N} | \dots \rangle \langle 6_{N-1} | \langle 6_{N-1} | \langle 6_{N-1} | \rangle \langle 6_{N-1} | \langle 6_{N-1} | \langle 6_{N-1} | \rangle \langle 6_{N-1} | \langle 6_{N-1} | \rangle \langle 6_{N-1} | \langle 6_{N-1} | \rangle \langle 6_{N-1} | \rangle \langle 6_{N-1} | \langle 6_{N-1} | \rangle \langle 6_{N-1} | \rangle \langle 6_{N-1} | \langle 6_{N-1} | \rangle \langle 6_{N-1} | \rangle \langle 6_{N-1} | \rangle \langle 6_{N-1} | \langle 6_{N-1} | \rangle \langle 6_{N$$

 \mathcal{B}_{β} is called right-normalized if it satisfies A three-leg terror

$$\mathcal{B}\mathcal{B}^{\dagger} = \mathbf{1} \quad \text{Explicitly:} \quad \left(\mathcal{B}\mathcal{B}^{\dagger}\right)_{\beta}^{\beta} = \mathcal{B}_{\beta}^{\delta\alpha} \mathcal{B}_{\alpha\delta}^{\dagger}^{\beta} = \mathbf{1}_{\beta}^{\beta} \tag{13}$$

Graphical notation for right-normalization: draw 'right-pointing diagonals' at vertices

When all B's are right-normalized, closing the zipper right-to-left is easy:

$$\langle \psi | \psi \rangle = \begin{pmatrix} x \\ x \end{pmatrix} = \begin{pmatrix} x$$

When all matrices of a MPS are right-normalized, the matrices for site N to any site $\ell = 1$

When all matrices of a MPS are right-normalized, the matrices for site N to any site $\ell = 1, ..., N$ define an orthonormal state space:

$$|\lambda\rangle = |\vec{\delta_{\ell}}\rangle [B^{\delta_{\ell}} B^{\delta_{\ell+1}} B^{\delta_{N}}]_{\lambda}^{\prime}$$
 (16)

$$\lambda'$$
 = λ' close the zipper

$$\langle \lambda' | \lambda \rangle = 1 \lambda'$$
 \otimes \otimes \otimes

Conclusion: MPS built purely from left-normalized A 's or purely from right-normalized B 's are automatically normalized to 1. Shorter MPSs built on subchains automatically define orthonormal state spaces.

Local operators act non-trivially only on a few sites (e.g. only one, or two nearest neighbors).

One-site operator

$$\hat{O}_{[\ell]} = |\sigma_{\ell}| > O_{\sigma_{\ell}} < \sigma_{\ell}|$$

$$\downarrow^{\sigma_{\ell}}$$

$$\downarrow^{\sigma_{\ell}}$$

$$\downarrow^{\sigma_{\ell}}$$

$$\downarrow^{\sigma_{\ell}}$$

$$\downarrow^{\sigma_{\ell}}$$

$$\downarrow^{\sigma_{\ell}}$$

E.g. for spin
$$\frac{1}{2}$$
: $(5_{\frac{1}{2}})^{\frac{1}{6}} = \frac{1}{2}(\frac{1}{2})$ $(5_{\frac{1}{2}})^{\frac{1}{6}} = (\frac{0}{10})$ $(5_{\frac{1}{2}})^{\frac{1}{6}} = (\frac{0}{10})$ (2)

In the Hilbert space of full system, one-site operator acts as unit operator on all sites except ℓ :

$$\hat{O}_{[\ell]} = |\vec{e}_{i}\rangle \underbrace{S_{e_{i}}^{e_{i}} \cdots S_{e_{i}}^{e_{i}}}_{O_{e_{i}} \cdots S_{e_{i}}} \underbrace{S_{e_{i}}^{e_{i}} \cdots S_{e_{i}}^{e_{i}}}_{O_{e_{i}} \cdots S_{e_{i}}} \underbrace{S_{e_{i}}^{e_{i}}}_{O_{e_{i}} \cdots S_{e_{i}}}}$$

Matrix element between two MPS:

$$\langle \tilde{\psi} \mid \hat{O}_{[\ell]} \mid \psi \rangle = \begin{cases} \langle \tilde{\psi} \mid \hat{O}_{[\ell]} \mid \psi \rangle \\ \langle \tilde{\psi} \mid \hat{O}_{[\ell]} \mid \psi \rangle \end{cases} = \begin{cases} \langle \tilde{\psi} \mid \hat{O}_{[\ell]} \mid \psi \rangle \\ \langle \tilde{\psi} \mid \hat{O}_{[\ell]} \mid \psi \rangle \end{cases} = \begin{cases} \langle \tilde{\psi} \mid \hat{O}_{[\ell]} \mid \psi \rangle \\ \langle \tilde{\psi} \mid \hat{O}_{[\ell]} \mid \psi \rangle \end{cases} = \begin{cases} \langle \tilde{\psi} \mid \hat{O}_{[\ell]} \mid \psi \rangle \\ \langle \tilde{\psi} \mid \hat{O}_{[\ell]} \mid \psi \rangle \\ \langle \tilde{\psi} \mid \hat{O}_{[\ell]} \mid \psi \rangle \end{cases} = \begin{cases} \langle \tilde{\psi} \mid \hat{O}_{[\ell]} \mid \psi \rangle \\ \langle \tilde{\psi} \mid \hat{O}_{[\ell]} \mid \psi \rangle \\ \langle \tilde{\psi} \mid \hat{O}_{[\ell]} \mid \psi \rangle \\ \langle \tilde{\psi} \mid \hat{O}_{[\ell]} \mid \psi \rangle \end{cases} = \begin{cases} \langle \tilde{\psi} \mid \hat{O}_{[\ell]} \mid \psi \rangle \\ \langle \tilde{\psi} \mid \hat{O}_{[\ell]} \mid \psi \rangle$$

The computation of such matrix element is simplest if $|\psi\rangle$ and $|\hat{\psi}\rangle$ are in 'site-canonical form', i.e. constructed from left- or right-normalized tensors for sites earlier or later than ℓ , respectively

site
$$\ell$$
 is special:

$$\alpha' \rightarrow \uparrow \stackrel{\epsilon_{c'}}{\longrightarrow} \beta'$$
 (6)

Matrix element:
$$\langle \widetilde{\psi} \mid \widehat{O} \mid \psi \rangle = \begin{cases} A & A & A & A & B & B & B \\ (\ell-1) & (7) &$$

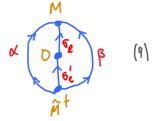
Close zipper from left using $C_{\ell^{-1}}$ from left-normalized β 's [see MPS-I.1-(15)], and from right using $D_{\ell^{+1}}$ from right-normalized B's [analogous to MPS-I.1-(20)].

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$$= \widetilde{M}_{\beta' \sigma_{\ell}' \alpha'}^{\dagger} C_{[\ell-1] \alpha}^{\alpha'} M^{\alpha \sigma_{\ell} \beta} D_{[\ell+1]_{\beta}}^{\alpha'} O_{\delta_{\ell}}^{\sigma_{\ell}'}$$
(8)

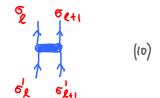
Now consider the expectation value, $\langle \psi | \hat{o} | \psi \rangle$ (i.e. drop all tilde's). The left-normalization of A''s guarantees that $C_{[\ell-1]} = 1$, and right-normalization of B''s that $D_{(\ell+1)} = 1$.

Hence



Two-site operator (e.g. for spin chain: \vec{S}_{ℓ} . $\vec{S}_{\ell+1}$)

$$\hat{O}_{[\ell,\ell+1]} = |\sigma_{\ell+1}^1\rangle|\sigma_{\ell}^1\rangle|O^{\sigma_{\ell}^1}|\sigma_{\ell+1}^1| < \sigma_{\ell}|< \sigma_{\ell+1}|$$



Matrix elements:

Matrix elements:

$$\langle \widetilde{\psi} \mid \widehat{O}_{\{\ell_i,\ell_{i+1}\}} \mid \psi \rangle = \langle \widetilde{\psi} \mid \widehat{O$$

$$= \widetilde{\mathcal{B}}_{\beta'\sigma'_{\ell+1}}^{\dagger} \widetilde{\mathcal{M}}_{\gamma'\sigma'_{\ell}}^{\dagger} \alpha' C_{[\ell-1]}^{\sigma'} \alpha \mathcal{M}^{\alpha\sigma_{\ell}\gamma} \mathcal{B}_{\gamma}^{\sigma_{\ell+1}} \mathcal{D}_{[\ell+2]_{\beta}}^{\sigma_{\ell+1}} \beta' \mathcal{D}_{\alpha'_{\ell}\sigma'_{\ell+1}}^{\sigma'_{\ell}\sigma'_{\ell+1}} (12)$$

Any matrix product can be expressed through different matrices without changing the product:

$$M M' = (M u u M') = \widetilde{M} \widetilde{M}'$$
 'gauge freedom

Gauge freedom can be exploited to 'reshape' MPSs into particularly convenient, 'canonical' forms:

Left-canonical (Ic-) MPS:

$$A^{\dagger}A = 1$$

(1)

Right-canonical (rc-) MPS:

Site-canonical (sc-) MPS:

$$M_{[\ell]}$$

$$= \times M_{[\ell]}^{\alpha \in \beta}$$

$$|\psi\rangle = |\vec{\epsilon}\rangle_{N} (A^{\epsilon_{1}} ... A^{\epsilon_{\ell-1}} M^{\epsilon_{\ell}}_{[\ell]} \mathcal{B}^{\epsilon_{\ell+1}} ... \mathcal{B}^{\epsilon_{N}}) = |\beta\rangle_{R} |\epsilon_{\ell}\rangle |\alpha\rangle_{L} M^{\alpha \epsilon_{\ell} \beta}$$

Bond-canonical (bc-) (or mixed) MPS:

$$|\psi\rangle = |\vec{\sigma}\rangle_{N}(A^{6!}...A^{6e})_{\alpha} S_{[R]\beta}^{\alpha\beta}(B^{6e+1}...B^{6N}) = \sum_{\alpha\beta} |\beta\rangle_{R} |\alpha\rangle_{L} S_{[R]}^{\alpha\beta}$$

$$(4)$$
can be chosen diagonal

How can we bring an arbitrary MPS into one of these forms?

Transforming to left-normalized form

14) = 18) (M61 ... M EN) Given:

[or with index: (SN) = X+1+1+1+> PM

Goal: left-normalize

$$\times \frac{A}{A} \stackrel{\text{M}}{\longrightarrow} \stackrel{\text{M}}{\longrightarrow} \stackrel{\text{M}}{\longrightarrow} \times \qquad (6)$$

Goal: left-normalize

Strategy: take a pair of adjacent tensors, MM^{\prime} , and use SVD,

$$MM' = USV^{\dagger}M' \equiv A\widetilde{M}$$
, with $A = U$, $\widetilde{M} = SV^{\dagger}M'$ (7)

$$\alpha \xrightarrow{M} M' \qquad SVD \propto \frac{1}{N} S \xrightarrow{N'} M' \qquad = \alpha \xrightarrow{M} \alpha' \qquad (8)$$

$$M^{\alpha \beta} M^{|\beta \delta'}_{\alpha'} = (U^{\alpha \delta}_{\alpha}) \langle S^{\lambda}_{\lambda'} V^{\dagger \lambda'}_{\beta} M^{|\beta \delta'}_{\alpha'} \rangle = A^{\alpha \delta}_{\alpha} \lambda \widehat{M}^{\lambda \delta'}_{\alpha'}$$
(9)

$$u^{\dagger}u = 1$$

 $u^{\dagger}u = 1$ ensures left-normalization: $A^{\dagger}A = 1$

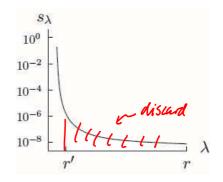
$$A^{\dagger}A = 1 \tag{(6)}$$

Truncation, if desired, can be performed by discarding some of

The smallest singular values,

$$\sum_{i=1}^{r}$$
 (but (10) remains valid!)

Note: instead of SVD, we could also me QR (cheaper!)



By iterating, starting from $M^{6_1}M^{6_2}$, we left-normalize

$$M^{61}$$
 to M^{6l-1} .

To left-normalize the entire MPS, choose $\ell \in \mathcal{N}$.

As last step, left-normalize last site using SVD on final $\, \stackrel{\scriptstyle \bullet}{\mathsf{M}} \,$:

$$\widetilde{M}^{\lambda \delta_{N}} = \underbrace{u^{\lambda \delta_{N}}}_{A^{\lambda \delta_{N}}} \underbrace{s_{1}^{\prime}}_{S_{1}^{\prime}} \underbrace{v_{1}^{\prime}}_{\delta_{N}^{\prime}} \underbrace{s_{1}^{\prime}}_{\delta_{N}^{\prime}} \underbrace{s_{1}^{\prime}}_{\delta_{N}^{\prime}$$

 $|\Psi\rangle = |\vec{\sigma}\rangle (A^{6}...A^{6N}) s$

diamond indicates single number

Ic-form:
$$(\psi) = (\vec{\sigma})_{N}(A^{6}, A^{6})_{N}$$

The final singular value, s_i determines normalization: $\langle \psi | \psi \rangle = |s_i|^2$.

Transforming to right-normalized form

Given:
$$(\gamma) = (\vec{r}) (M^6 ... M^6)$$

16) (Mar... 191 a) 4 4 4 4 6

Goal : right-normalize M^{6N} to M^{6Q+1}

Strategy: take a pair of adjacent tensors, MM , and use SVD:

$$MM' = MUSU^{\dagger} \equiv \widetilde{M}B$$
 with $\widetilde{M} = M^{\dagger}US$ $B = V^{\dagger}$. (13)

$$M_{\alpha}^{\ \ \sigma\beta} M_{\beta}^{\ \ \sigma'\alpha'} = \left(M_{\alpha}^{\ \ \sigma\beta} U_{\beta}^{\ \ \lambda} S_{\lambda}^{\lambda'}\right) \left(V_{\lambda'}^{\dagger} \sigma'\alpha'\right) = \widetilde{M}_{\alpha}^{\ \ \sigma\lambda'} B_{\lambda'}^{\ \ \sigma'\alpha'}$$
 (15)

Here, $V^{\dagger}V = 1$ ensures right-normalization: $B^{\dagger} = 1$. (6)

Starting form $M^{6N-1}M^{6N}$, move leftward up to $M^{6N}M^{6N+1}$

To right-normalize entire chain, choose / and at last site, $\ell = 1$

$$\widetilde{M}_{1}^{G,\lambda} = \underbrace{U_{1}^{G,\lambda} \underbrace{S_{1}^{G,\lambda} \underbrace{V_{1}^{G,\lambda}}_{B_{1}^{G,\lambda}}}$$
 . So determines normalization. (13)

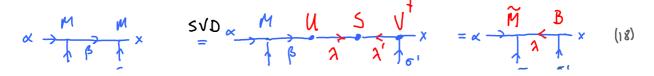
Exercise

(a) Right-normalize a state with right-pointing arrows!



Hint: start at

and note the up \(\leftrightarrow\) down changes in index placement.



(b) Left-normalize a state with left-pointing arrows!

Hint: start at M⁶¹M⁶²:

Transforming to site-canonical form

Then right-normalize sites N to ℓ , starting from site N .

Result:

$$|\psi\rangle = |\delta_{N}\rangle ... |\delta_{R^{+}}\rangle (B^{\delta_{R^{+}}}...B^{\delta_{N}}) |\delta_{R}\rangle |\delta_{R^{-}}\rangle ... |\delta_{L}\rangle (A^{\delta_{L}}...A^{\delta_{R^{-}}}) |\alpha\rangle (23)$$

$$= |\beta\rangle_{R} |\delta_{R}\rangle |\alpha\rangle_{L} |\Delta_{L}\rangle (24)$$

The states $\langle \alpha, \sigma_{\ell}, \beta \rangle \equiv \langle \beta \rangle_{\ell} \langle \sigma_{\ell} \rangle \langle \alpha \rangle_{\ell}$ form an orthonormal set:

$$\langle \alpha', \sigma'_{\ell}, \beta' | \alpha, \sigma_{\ell}, \beta \rangle = \delta^{\alpha'}_{\alpha} \delta^{\sigma_{\ell}'}_{\sigma_{\ell}} \delta^{\beta'}_{\beta}$$
 (25)

(Exercise: verify this, using $A^{\dagger}A = 1$ and $BB^{\dagger} = 1$.)

Transforming to bond-canonical form

Transforming to bond-canonical form

$$\overline{M} = U S V^{\dagger}$$
 $A = U \hat{S} = V^{\dagger} B$ (Exercise: add indices!) (27)

The states

$$\left(\frac{\lambda}{\lambda}, \frac{\lambda'}{\lambda'}\right) = \left(\frac{\lambda}{\lambda'}\right)$$
 form an orthonormal set.

$$\langle \bar{\lambda}, \bar{\lambda}' | \lambda, \lambda' \rangle = \delta^{\bar{1}}_{\lambda} \delta^{\lambda'}_{\lambda}$$
 (29)

This is called the 'local bond basis for bond ℓ ' (from site ℓ to ℓ). It has dimension ℓ . (ℓ = dimension of singular matrix ℓ).

 $\left(\begin{array}{c} \lambda \\ \lambda \end{array}\right) = \left(\begin{array}{c} \lambda^{1} \\ \lambda \end{array}\right)_{R} \left(\begin{array}{c} \lambda \\ \lambda \end{array}\right)_{L}$ form 'local bond basis' for bond ℓ -1 (from site ℓ -1 to ℓ).