Sym-III.1

For representation theory of SU(N), see [Alex2011]

### 1. Motivation

'Symmetries II: non-Abelian' showed: in the presence of symmetries, A-tensors factorize:

$$A^{(0,i;q),(R,j;r)}(S,k;s) = (\hat{A}^{0,R}S)^{ij}_{k}(C^{0,R}S)^{ij}_{s}$$

$$Q_{i,i;q}A_{i,j;s} = Q_{i,i}A_{i,j}S_{i,j}$$

$$Q_{i,i;q}A_{i,j;s} = Q_{i,i}A_{i,j}S_{i,j}S_{i,s}$$

$$Q_{i,i;q}A_{i,j;s}S_{i,j}S_{i,s}$$

Goal: reduce explicit reliance on Clebsch-Gordan tensors (CGT) as much as possible!

Why? CGT can become very large objects for groups of large rank, e.g. SU(N) with N > 3. Hence, whenever possible, avoid computing and contracting them explicitly.

## **Multiplet dimensions**

Irreducible representation (irrep) of symmetry group forms a vector space:

$$V^{Q} = span \{Q, q\}, \qquad , \qquad q = 1, ..., dim(V^{Q}) = d_{Q} = |Q|$$
'irrep' label or 'internal' label, distinguishes states in multiplet. (1)

In general, internal label is a composite label,  $q = (q_1, q_2, \dots, q_r)$  where  $\tau$  is the 'rank' of the group, i.e. the number of commuting generators,  $\hat{S}_1, \dots, \hat{S}_r$ 

These span the 'Cartan subalgebra', can be diagonalized simultaneously:

$$\hat{S}_{i}|_{q}\rangle = q_{i}|_{q}\rangle , \quad i=1,..., f$$
 (3)

Their eigenvalues can be used to label a basis for  $\c^6$  , as done in (1).

Multiplet dimension: 
$$|Q| = \prod_{i=1}^{T} q_{i}^{max} \sim \left(q_{averag}^{max}\right)^{T}$$
 (4)

For SU(N), 
$$\uparrow = V - 1$$
, typical dimensions grow as  $\sim 10^{1/2}$ , 'large' for  $10^{1/2}$  (s)

Hence: efficient numerics tries to avoid working 'inside' multiplets; rather treat them as closed units.

2. Outer multiplicity Sym-III.2

Decomposition of tensor product of two irreps into direct sum of irreps:

$$\bigvee^{Q} \otimes \bigvee^{Q'} = \sum_{i} \bigvee^{QQ'} \bigvee^{Q''} \bigvee^{Q''} = \sum_{i} \sum_{j} \bigvee^{QQ''} \bigvee^{Q''} \bigvee^{Q'} \bigvee^{Q''} \bigvee^{Q'} \bigvee^{Q}$$

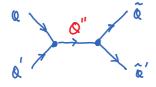
'Outer multiplicity' (OM)  $\bigvee^{QQ'}_{Q''}$  is an integer specifying how often the irrep  $\bigvee^{Q}_{Q'}$  occurs in the decomposition of the direct product  $\bigvee^{Q}_{Q'}$ .

For given Q'', 'outer multiplicity index'  $p_{ij} = 1, \dots, p_{ij} = 1, \dots, p_{ij$ 

For other groups, e.g. 
$$SU(N \ge 3)$$
, the OM can be  $> 1$ .

The extra OM-index brings additional complexity to tensor network codes. 'SU(3) is much harder than SU(2)'.

But also for SU(2), OM enters when coupling more than two irreps:



For given irreps  $\delta$ , Q',  $\widetilde{Q}$ ,  $\widetilde{Q}'$ , there can be several possible choices for Q''. The total number of possibilities,  $\mu$  is the OM of Q''.

## Clebsch-Gordan tensors (CGT)

transforms generators into block-diagonal form: (drawn for N = 3, N = 3

The basis transformation ( is encoded in Clebsch-Gordan tensors (CGTs):

$$|Q''_{\mu,q} : Q . Q' \rangle = \sum_{q'} |Q'_{q'} : Q |Q |Q | Q'_{q'} |Q'_{\mu q'} : Q |Q' | Q'_{\mu q'} : Q |Q'_{\mu q'} : Q$$

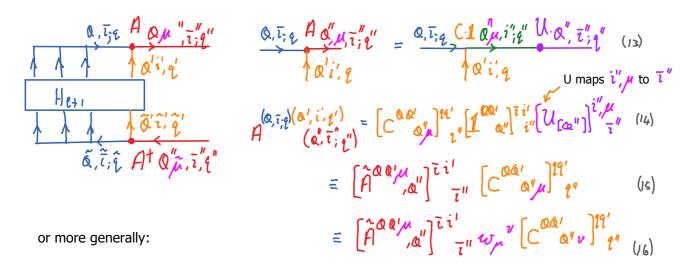
Rank-3 CGTs are sometimes called '3-j symbols', since they link 3 irreps.

Factorization of A-tensor (see Sym-II.15) must account for OM:

Recall iterative diagonalization, unitary transformation into energy eigenbasis:

Hrong Va, collective index: direct product structure (12)

Combined transformation from old energy eigenbasis to new energy eigenbasis:



A-matrix factorizes, into product of reduced A-matrix and CGT!!

A 'does not know' about OM. But its factorization does! This structure can be exploited to reduce numerical costs: Associate with A rather than with C

 $A = (\widetilde{A}^{\mu} w_{\mu}) C_{\nu} = \widetilde{\widetilde{A}}^{\nu} C_{\nu}$ 

3. Arrow inversion Sym-III.3

CGTs can always be chosen real. According to (13), they represent a unitary transformation. Hence, for fixed Q, Q', they satisfy:

$$\sum_{\mathfrak{d}\mathfrak{d}'} \left( c^{+\tilde{\mathfrak{d}}'',\tilde{\mathfrak{p}}}_{\mathfrak{d}'\mathfrak{d}} \right)^{\tilde{\mathfrak{l}}''} _{\mathfrak{l}'\mathfrak{l}} \left( c^{\mathfrak{d}\mathfrak{d}'}_{\mathfrak{d}'',\tilde{\mathfrak{p}}} \right)^{\mathfrak{l}\mathfrak{l}'} _{\mathfrak{l}''} = \mathbf{1}^{\tilde{\mathfrak{d}}''}_{\mathfrak{d}''} \mathbf{1}^{\tilde{\mathfrak{l}}''}_{\mathfrak{l}''} \mathbf{1}^{\tilde{\mathfrak{l}}''}_{\mathfrak{l}''}$$

$$(14)$$

$$\sum_{Q''_{i,j}} (C^{\alpha \alpha'_{i,j}})^{\widetilde{l}\widetilde{l}'_{j,j}} (C^{Q'_{i,j}}, \alpha' \alpha)^{2''_{i,j}} = 1^{\frac{2}{7}} 1^{\frac{2}{7}}, \qquad (15)$$

Note: 
$$Z (s) \Big|_{\widetilde{q}=q} = I_{q'}^{\widetilde{q}'} |Q|$$
  $= \int |Q'|$   $= \int |Q'|$ 

Weichselbaum2019 uses a different normalization, such that, for fixed Q,Q',Q",

'full contraction of all indices except OM index' yields:

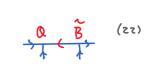
Then, 'opening' any leg yields unit matrix divided by dimension of that leg:

$$Q'' = \frac{1}{|Q'|} = \frac{1}{|Q'|} = \frac{1}{|Q'|} = \frac{1}{|Q''|} = \frac{1}{|Q''|}$$

Prefactor on r.h.s. follows from requirement that trace over open leg reproduces (20).

#### Inverting arrows

Recall general procedure:



(17)

(18)

How does one invert arrows in CGT-sector?

Define 
$$\left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/2} = \prod_{Q \in Q} \left( \bigcup_{Q \in Q} Q \right)^{1/$$

and its conjugate irrep Q to the trivial, one-dimensional irrep *O* .

dimensions: 
$$|Q| = |\overline{Q}|$$
  $|D| = 1$   $(24)$ 

Then U is unitary: 
$$u^{Q\bar{c}} \circ u^{\dagger} \circ Q = 1^{|Q|}$$
, (254)  $u^{\dagger} \circ Q \circ U^{Q\bar{c}} \circ Q = 1^{|Q|}$  (255)

Graphical argument shows why: consider

$$T_{r}$$
  $V_{qq}$   $V_$ 

now open Q-leg:

now open Q-leg:
$$u \overset{\circ}{\circ} u \overset{\dagger}{\circ} o = 0 \qquad u \overset{\circ}{\circ} u \overset{\circ}{\circ} u = 1 \overset{\circ}{\circ} u \overset{\circ}{\circ} u$$

Similarly, opening (Leg leads to (25b). Compact graphical notation: drop dashed loop

hence arrows can be inverted by inserting  $uu^{\dagger} = 1$  or  $u^{\dagger}u = 1$ (26)

U is sometimes called '1-j symbol', since it involves only a single irrep and its conjugate.

U can be computed by finding the ground state of the pseudo-Hamiltonian acting on  $\bigvee^{a}$ 

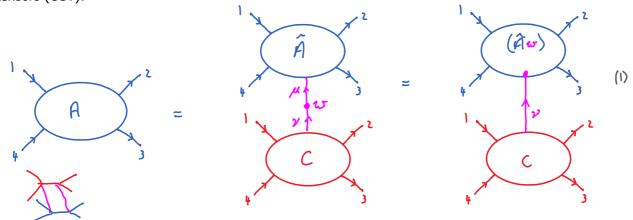
$$\hat{H}_{Q} = \sum_{\alpha=1}^{T} \hat{S}_{\alpha} \hat{J}_{\alpha} + \hat{S}_{\alpha} \hat{S}_{\alpha}^{\dagger}$$
(78)

Let 
$$|0\rangle = |0, q\rangle u^{q}$$
,  $|0\rangle = |0, q\rangle u^{q}$ , then  $|0\rangle = |0, q\rangle u^{q}$  since this maps  $|0\rangle = |0, q\rangle u^{q}$ 

# 4. Pairwise contractions and X-symbols

Sym-III.4

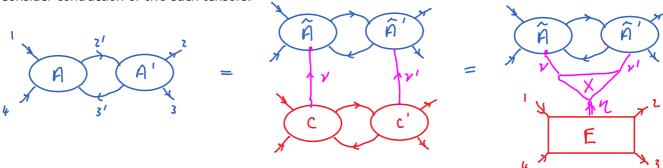
Consider a four-leg tensor, factorized into reduced matrix element tensors (RMT) and Clebsch-Gordan tensors (CGT):

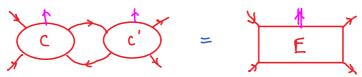


The OM-matrix  $\, \boldsymbol{w} \,$  can always be contracted onto the RMT, as indicated on the right.

 $\hat{\rho}$  is in active memory (has to be stored, updated, etc.), whereas is 'known' (stored in library on hard disk).

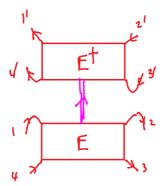
Consider contraction of two such tensors:

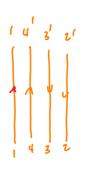




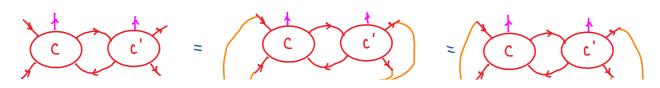
Resolve identity on space of all open legs:

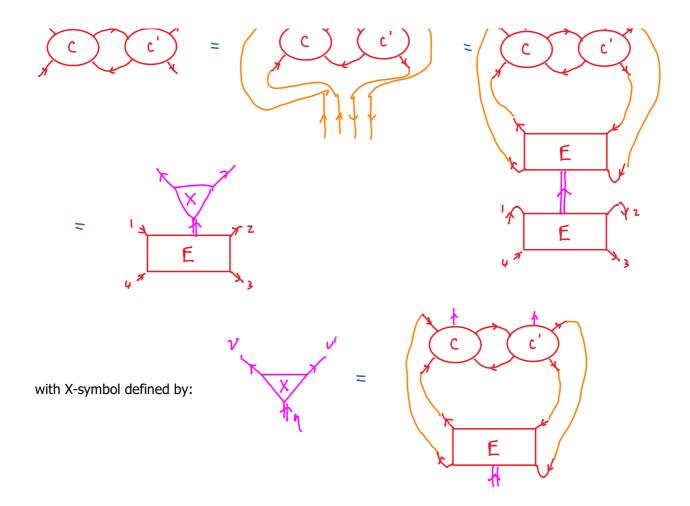
$$E^{+}E = 1$$
:





Then we obtain:





Manipulations with A's happen in active memory, those with Cs, Es, Xs are done only once, then stored on hard disk for to be contracted in active memory. This brings huge reduction in numerical costs, since Cs, Es can be huge objects, whereas Xs are small.