< 4/4> 1. Overlaps and normalization

Consider overlap of 2-site MPS:

$$|\beta\rangle = |\varepsilon_2\rangle|\varepsilon_1\rangle A^{(\varepsilon)} \alpha B^{d\varepsilon_1} \beta \qquad (1)$$

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$$\langle \beta^{(1)}\beta^{(2)} = A^{16_{1}} B^{d6_{1}} B^{d6_{1}} \beta^{(2)} \langle G_{1}^{(1)}|\langle G_{2}^{(1)}|G_{1}\rangle A^{16_{1}} B^{d6_{2}} \beta^{(2)}$$
(2)

introduce
$$A^{\dagger}$$
's: $A^{\dagger} \alpha' \varsigma_{1} B^{\dagger} \delta_{2} \alpha^{\dagger}$ (3)

reorder
$$A^{\dagger}$$
's: $B^{\dagger} \beta_{\sigma_{2}}^{\prime} A^{\dagger} A^{\prime} \sigma_{1}^{\prime} A^{\prime} \sigma_{2}^{\prime} B^{\prime} \beta_{2} \beta_{2}$

Use diagrammatic rules to keep track of contraction patterns:

Ket:
$$[\beta] = [\sigma] \alpha \beta A^{\alpha \sigma} \beta$$
 (5) $\alpha \beta^{\alpha \sigma} \beta$
Bra: $\langle \beta | = \langle \alpha | \langle \sigma | A^{\alpha \sigma} \beta = A^{\dagger} \beta_{\sigma \alpha} \langle \alpha | \langle \sigma | (6) \alpha \beta^{\alpha \sigma} \beta$

We accommodated complex conjugation via Hermitian conjugation and index transposition:

At BER = AxEB

This moves upstairs indices downstairs and vice versa, i.e. invents all arrows in diagram. Note that in diagram vertex, α sits left, β right, whereas on A^{\uparrow} , β sits left, α right. This convention may seem initially awkward, but it greatly simplifies the structure of diagrams representing overlaps.

Generalization to many-site MPS:

$$| \psi \rangle = | \overline{e} \rangle A_{[1]}^{1} \alpha A_{[2]}^{\alpha} \beta A_{[3]}^{\beta} \gamma \dots A_{[N]}^{\alpha} \gamma \dots A_{[N]}^{\alpha}$$

$$(z)$$

$$dummy index$$

Square brackets indicate that each site has a different *A* matrix. We will often omit them and use the shorthand, $A \stackrel{\alpha \circ \beta_{\ell}}{=} \beta = A \stackrel{\alpha \circ \beta_{\ell}}{[\ell]} \beta$ since the ℓ on σ_{ℓ} uniquely identifies the site.

dummy site

Bra:

$$A_{[1]} A_{[2]} A_{[2]} A_{[3]} A_{[$$

MPS-I.1

x.



Recipe for ket formula: as chain grows, attach new matrices on the right (in <u>same</u> order as vertices in diagram); resulting in a matrix product structure.

Bra:

$$\begin{cases} \zeta_{1}|\zeta_{5}|...,\zeta_{N}| \\ \zeta_{V}| = & \zeta_{0}^{5}| \overline{A}_{(1)}^{1_{G_{1}}} \overline{A}_{(2)}^{\alpha_{5}} \overline{A}_{(3)}^{\beta_{5}} \overline{Y} \cdots \overline{A}_{(N)}^{\mu_{5}} = & t_{1} \\ N \\ \zeta_{V}| = & A_{(1)}^{1_{G_{1}}} \overline{A}_{(1)}^{1_{G_{1}}} \overline{A}_{(2)}^{\alpha_{5}} \overline{A}_{(3)}^{\beta_{5}} \overline{Y} \cdots \overline{A}_{(N)}^{\mu_{5}} = & A_{(N)}^{1_{G_{1}}} \overline{\sigma}_{3\beta} \overline{A}_{(2)}^{1_{G_{1}}} \overline{\sigma}_{2\alpha} \overline{A}_{(1)}^{1_{G_{1}}} \overline{\sigma}_{1} \overline{\zeta_{0}} = \\ & X \\ X \\ A_{(1)}^{1_{G_{1}}} \overline{A}_{(2)}^{1_{G_{1}}} \overline{A}_{(3)}^{1_{G_{1}}} \cdots \overline{A}_{(N)}^{n_{N}} = & X \\ A_{(1)}^{1_{G_{1}}} \overline{A}_{(2)}^{1_{G_{1}}} \overline{A}_{(3)}^{1_{G_{1}}} \cdots \overline{A}_{(N)}^{n_{N}} = \\ & X \\ A_{(1)}^{1_{G_{1}}} \overline{A}_{(2)}^{1_{G_{1}}} \overline{A}_{(3)}^{1_{G_{1}}} \cdots \overline{A}_{(N)}^{n_{N}} = & X \\ & X$$

We expressed all matrices via their Hermitian conjugates by transposing indices and inverting arrows. To recover a matrix product structure, we ordered the Hermitian conjugate matrices to appear in the <u>opposite</u> order as the vertices in the diagram.

Recipe for bra formula: as chain grows, attach new matrices \mathbf{A}^{\dagger} on the left, opposite to vertex order in diagram.

Now consider overlap between two MPS:

$$\langle \hat{\psi} | \psi \rangle = \begin{pmatrix} A_{(1)} & A_{(1)} & A_{(3)} & A_{(3)} \\ & & & & \\ & & & \\ & & & \\ & & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\$$

$$= \widetilde{A}_{(w)} \widetilde{\epsilon}_{N} \mu' \cdots \widetilde{A}_{(3)} \widetilde{\epsilon}_{N} \mu' \widetilde{A}_{(2)} \widetilde{\epsilon}_{z\alpha'} \widetilde{A}_{(1)} \widetilde{\epsilon}_{(1)} \widetilde{\epsilon$$

Exercise: derive this result algebraically from (7), (9),

If we would perform the matrix multiplication first, for fixed \vec{r} , and then sum over \vec{r} , we would get d^N terms, each of which is a product of 2N matrices. Exponentially costly!

Diagrammatic depiction: 'closing zipper' from left to right.

$$C_{[a]} \xrightarrow{\alpha} \beta \xrightarrow{\beta} \beta$$

The set of two-leg tensors $C_{[\ell]}$ can be computed iteratively:

Initialization:

$$\zeta_{[o]} \begin{pmatrix} \uparrow k \\ \leftarrow k \end{pmatrix} = \begin{pmatrix} \\ (identity) \end{pmatrix}$$



 $C'_{[o]} = 1$

Iteration step:

$$C_{[\ell]} = C_{[\ell-i]} = C_{[\ell-i]}$$

Final answer:

 $\langle \tilde{\psi} | \psi \rangle = C'_{[N]}$



(14)

Cost estimate (if all A's are p_{γ} D):

Remark: a similar iteration scheme can be used to 'close zipper from right to left':



When all A's are left-normalized, closing the zipper left-to-right is easy, since all c_{ℓ} reduce to identity matrices:

$$C_{[o]} = \begin{bmatrix} 1 \\ 1 \end{bmatrix} C_{[i]_{a}}^{a'} = \begin{bmatrix} 1 \\ a' \end{bmatrix} C_{[i]_{a}}^{a'} = \begin{bmatrix} 1 \\ a' \end{bmatrix} C_{[e]}^{a'} = \begin{bmatrix} 1 \\ 1 \end{bmatrix} C_{[e]}^{a'} =$$

identity mutix

Hence:

$$\langle 24|2\rangle =$$
 (24)

When all matrices of a MPS are left-normalized, the matrices for site 1 to any site $\ell = 1, ..., N$ define an orthonormal state space:



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$$\langle \chi' | \chi \gamma = I \chi' \chi$$

So far we have viewed an MPS as being built up from left to right, hence used right-pointing arrows on ket diagram. Sometimes it is useful to build it up from right to left, running left-pointing arrows.

Building blocks:

Iterating this, we obtain kets and bras of the form

$$|\psi\rangle = |f_N\rangle|f_{N-r}\rangle \dots |f_1\rangle B_1^{f_1\lambda} \dots B_k^{f_{N-r}\lambda} B_{\alpha}^{f_{N-r}}$$

$$\langle \mathfrak{P} | = \mathfrak{B}_{16_{N}}^{\dagger} \mathfrak{B}_{\alpha \mathfrak{E}_{N-1}}^{\dagger} \mathfrak{B}_{\alpha \mathfrak{E}_{N-1}}^{\dagger} \mathfrak{B}_{\beta \mathfrak{E}_{1}}^{\dagger} \langle \mathfrak{E}_{1} | \ldots \langle \mathfrak{E}_{N-1} | \langle \mathfrak{E}_{N} \rangle$$

is called right-normalized if it satisfies Bß A three-leg terror (BB^{\dagger}) = $B^{\sigma \alpha} B^{\dagger}$ B^{\prime} $BB^{\dagger} = 1$

Explicitly:
$$(BB')$$

(31)

(29)

(36)

Graphical notation for right-normalization: draw 'right-pointing diagonals' at vertices B



When all B's are right-normalized, closing the zipper right-to-left is easy:

$$\langle \psi|\psi\rangle = \frac{x}{x} + \frac{y}{x} + \frac{y}{$$

When all matrices of a MPS are right-normalized, the matrices for site N to any site $\ell = 1, ..., N$ define an orthonormal state space:



Conclusion: MPS built purely from left-normalized A 's or purely from right-normalized B's are automatically normalized to 1. Shorter MPSs built on subchains automatically define orthonormal state spaces.

Any matrix product can be expressed through different matrices without changing the product:

$$M M' = (M u) u M') = \widetilde{M} \widetilde{m}'$$
 'gauge freedom'

Gauge freedom can be exploited to 'reshape' MPSs into particularly convenient, 'canonical' forms:

How can we bring an arbitrary MPS into one of these forms?

14) = (=) (M " ... M ")

M⁶¹ to M⁶e-(

 $|S_N\rangle = x^{1} + 1 + 1 + 2N$

Transforming to left-normalized form

Given:

[or with index:

Goal : left-normalize

Goal : left-normalize

to M^Ge-(M 61

 $|\psi\rangle = (\overline{c})_{N}(A^{6}, ..., A^{6N}) s_{1}$

MM' , and use SVD, Strategy: take a pair of adjacent tensors,

$$MM' = USV^{\dagger}M' = A\widetilde{M}, \text{ with } A = U, \widetilde{M} = SV^{\dagger}M'$$
 (7)

$$\alpha \xrightarrow{\mathcal{M}} \underset{\sigma}{\overset{\mathcal{M}}{\longrightarrow}} \underset{\sigma}{\overset{\mathcal{M}}{\longrightarrow}} \alpha' \xrightarrow{S \sqrt{D}} \underset{\sigma}{\overset{\mathcal{N}}{\longrightarrow}} \overset{\mathcal{N}}{\overset{\mathcal{N}}{\longrightarrow}} \overset{\mathcal{N}}{\overset{\mathcal{N}}{\longrightarrow}} \overset{\mathcal{N}}{\overset{\mathcal{N}}{\longrightarrow}} \alpha' = \alpha \xrightarrow{\mathcal{A}} \overset{\mathcal{M}}{\overset{\mathcal{M}}{\longrightarrow}} \alpha' \quad (8)$$

$$M^{\alpha \sigma} \beta M^{\beta \sigma'} \alpha' = (U^{\alpha \sigma} \lambda) (s^{\lambda} \lambda' V^{\dagger} \lambda' \beta M^{\beta \sigma'} \alpha') = A^{\alpha \sigma} \lambda \widetilde{M}^{\lambda \sigma'} \alpha' \qquad (9)$$

The properly

lc-form:

 $\mathcal{U}^{\dagger}\mathcal{U} = \mathbf{1}$ ensures left-normalization:

$$A^{\dagger}A = \mathbf{1} \qquad (10)$$

Truncation, if desired, can be performed by discarding some of

The smallest singular values,

$$\sum_{\substack{n=1\\n}}^{n} \sum_{\substack{n=1\\n}}^{n} \sum_{\substack{n=1\\n}}^{n} (but (10) remains valid!)$$

Note: instead of SVD, we could also me QR (cheaper!)

By iterating, starting from $M^{6_1} M^{6_2}$, we left-normalize

To left-normalize the <u>entire</u> MPS, choose $\mathcal{L} = \mathcal{N}$. As last step, left-normalize last site using SVD on final $\,\widetilde{\,eaglefontume}\,$:

$$\widetilde{M}^{\lambda 6_{\mathcal{N}}} = \underbrace{u^{\lambda 6_{\mathcal{N}}}}_{A^{\lambda 6_{\mathcal{N}}}, S_{1}} \underbrace{v^{\dagger}}_{S_{1}} \underbrace{\lambda^{\dagger}}_{S_{1}} = \underbrace{u S v^{\dagger}}_{S_{1}} \underbrace{v^{\dagger}}_{S_{1}} \underbrace{\lambda^{\dagger}}_{S_{1}} = \underbrace{u S v^{\dagger}}_{S_{1}} \underbrace{v^{\dagger}}_{S_{1}} \underbrace{v^{\dagger}}_{S_{1}$$

diamond indicates single number







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M61

$$\times \frac{H}{5} \xrightarrow{H} f \xrightarrow{H}$$

Ic-form:
$$(\psi) = (\vec{\sigma})_{\nu} (A^{6}' \dots A^{6} N)_{S_{1}}$$

עומודוטרוע ווועוכמנכס single number

(16)

The final singular value, *S*, determines normalization:

$$\langle \psi | \psi \rangle = | \varsigma_i |^2$$
. (12)

Transforming to right-normalized form

 $|\psi\rangle = |\vec{e}\rangle_{I} (M^{e_{I}} \dots M^{e_{N}})$ Given: [or with index: $|S_1\rangle \approx S_1 \leftarrow 1 \leftarrow 1 \leftarrow 1$ ٦

Goal : right-normalize M^{ℓ_N} to $M^{\ell_{\ell+1}}$





Strategy: take a pair of adjacent tensors, MM' , and use SVD:

 $MM' = MUSU^{\dagger} = \widetilde{M}B$ with $\widetilde{M} = MUS$, $B = U^{\dagger}$. (13) $\alpha \xrightarrow{M} \underset{\beta}{\overset{M}{\longrightarrow}} \overset{M}{\xrightarrow{\beta}} \overset{\beta}{\xrightarrow{\gamma}} \overset{\beta}{\xrightarrow{\gamma}}$

$$M_{\alpha}^{\beta\beta} M_{\beta}^{\delta'\alpha'} = (M_{\alpha}^{\beta\beta} U_{\beta}^{\lambda} S_{\lambda}^{\lambda'}) (V_{\lambda'}^{\dagger \sigma'\alpha'}) = \widetilde{M}_{\alpha}^{\delta\lambda'} B_{\lambda'}^{\delta'\alpha'} (S)$$

Here, $\sqrt{\frac{1}{2}} \sqrt{\frac{1}{2}} = 1$ ensures right-normalization: $\frac{1}{2} \frac{1}{2} = 1$.

Starting form $\mathcal{M}^{\mathfrak{G}_{\mathcal{N}-1}} \mathcal{M}^{\mathfrak{G}_{\mathcal{N}}}$, move leftward up to $\mathcal{M}^{\mathfrak{G}_{\mathcal{R}}} \mathcal{M}^{\mathfrak{G}_{\mathcal{R}+1}}$

To right-normalize entire chain, choose / and at last site, $\lambda = 1$

 $\widetilde{\mathcal{M}}_{1}^{\sigma,\lambda} = \underbrace{\mathcal{U}_{1}^{\prime}S_{1}^{\prime}V_{1}^{\dagger}G_{1}\lambda}_{=, S_{1}^{\prime}S_{1}^{\prime}} \underbrace{\mathcal{J}_{1}^{\dagger}G_{1}\lambda}_{B_{1}^{\sigma,\lambda}} . \qquad S_{1}^{\prime} \text{ determines normalization.}$ (17)

Exercise

(a) Right-normalize a state with right-pointing arrows!

Hint: start at

and note the up <-> down changes in index placement.

MGN-1 MGN



$$\langle \alpha', \sigma_{\ell}', \beta' \mid \alpha, \sigma_{\ell}, \beta \rangle = \delta_{\alpha}^{\alpha'} \delta_{\delta_{\ell}}^{\sigma_{\ell}'} \delta_{\beta}^{\beta'}$$
 (25)

(Exercise: verify this, using $A \stackrel{t}{A} = 1$ and $B \stackrel{t}{B} \stackrel{t}{=} 1$.) This is 'local site basis' for site ℓ . Its dimension $\mathcal{D}_{a} \cdot d \cdot \mathcal{D}_{\beta}$ is usually $\langle \langle \langle d \rangle \rangle$ of full Hilbert space.

Transforming to bond-canonical form

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Transforming to bond-canonical form

Start from (e.g.) sc-form, use SVD for $\overline{M} \approx U S V^{\dagger}$, combine $\mathbb{O} \quad V^{\dagger}$ with neighboring \mathcal{B} , or $\mathbb{O} \quad \mathcal{V}$ with neighboring \mathcal{A} .

$$\langle \bar{\lambda}, \bar{\lambda}' | \lambda, \lambda' \rangle = \langle \bar{\lambda}, \bar{\lambda}' \rangle$$
(25)

This is called the 'local bond basis for bond ℓ ' (from site ℓ to ℓ). It has dimension ℓ . (τ = dimension of singular matrix S).

 $\left(\lambda,\lambda'\right) = \left(\lambda'\right)_{R} \left(\lambda\right)_{L}$ form 'local bond basis' for bond $\ell - 1$ (from site $\ell - 1$ to ℓ).